



# Decentralised storage in the blockchain era: solutions and challenges

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**CryptoNet**



**Protocol Labs**

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Torino

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Filecoin Protocol upgrades

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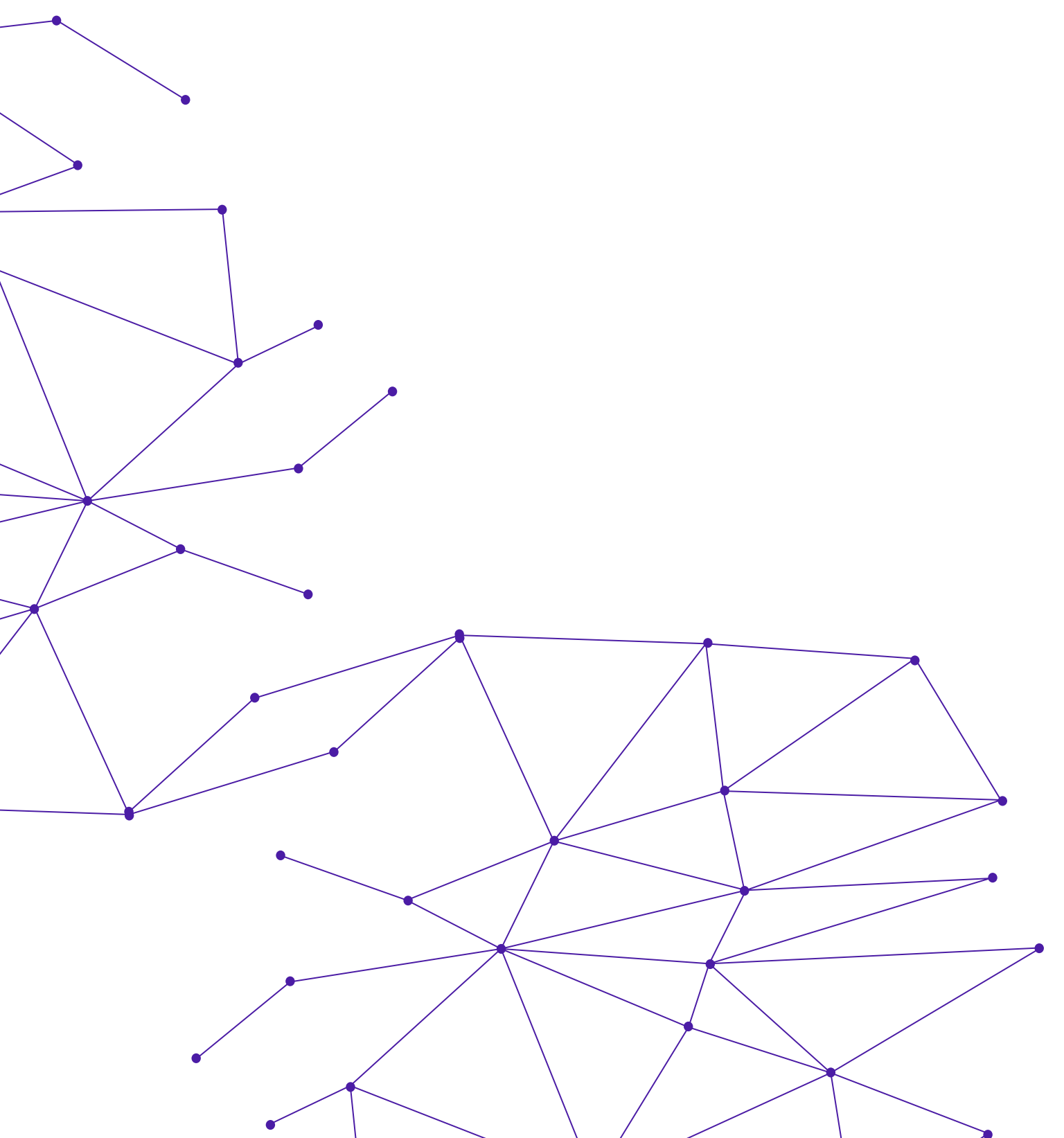
Apply for grants:

 [Grants \[Paused\]](#)

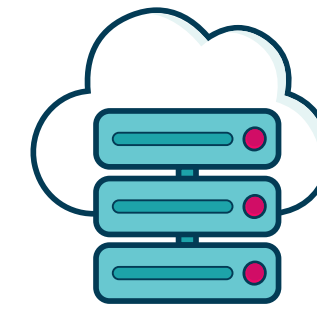


# Talk Outline

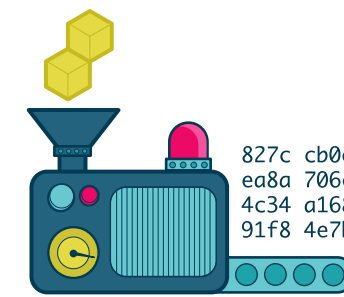
## Three Topics:



1. Decentralised storage



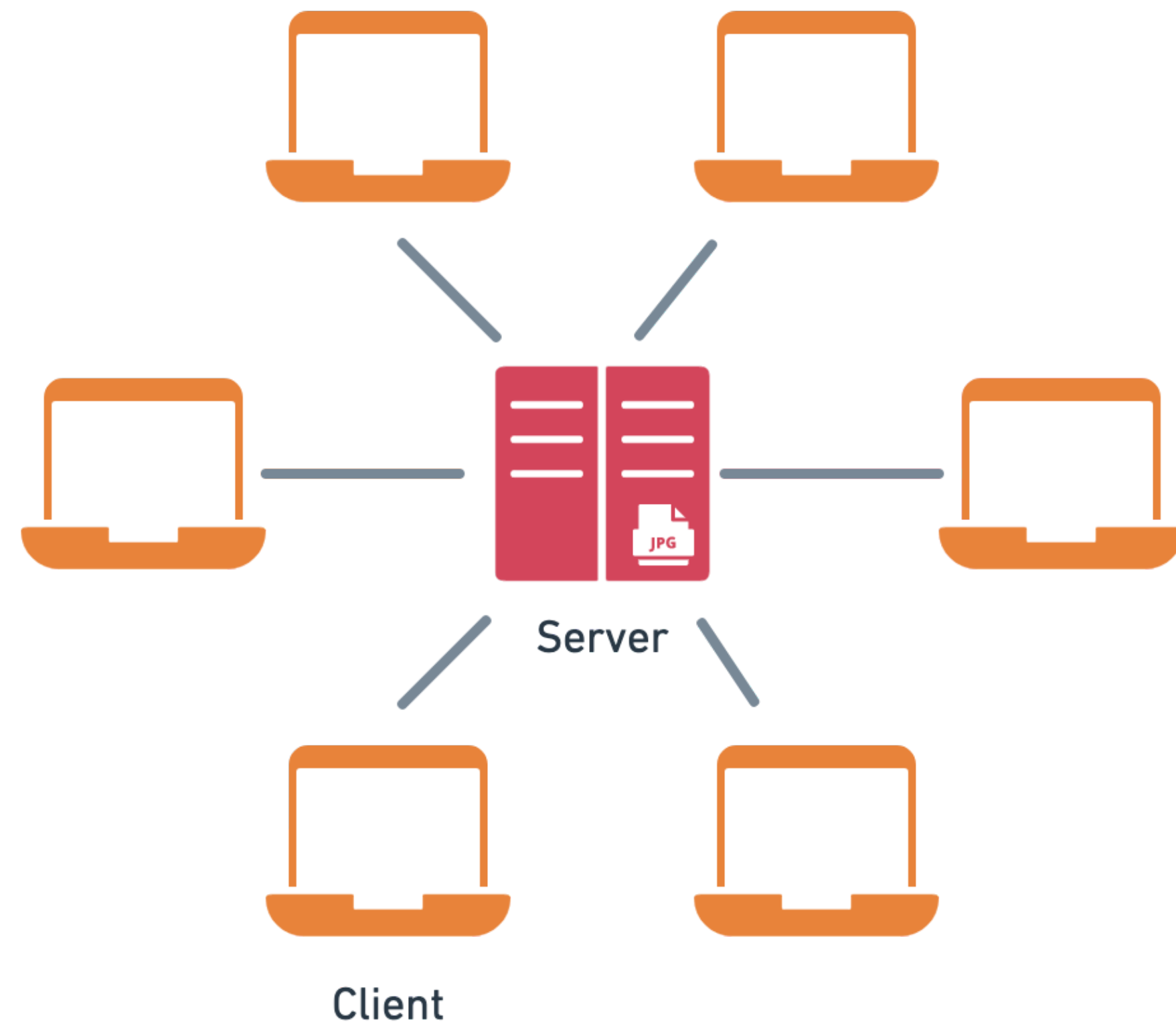
2. Blockchain



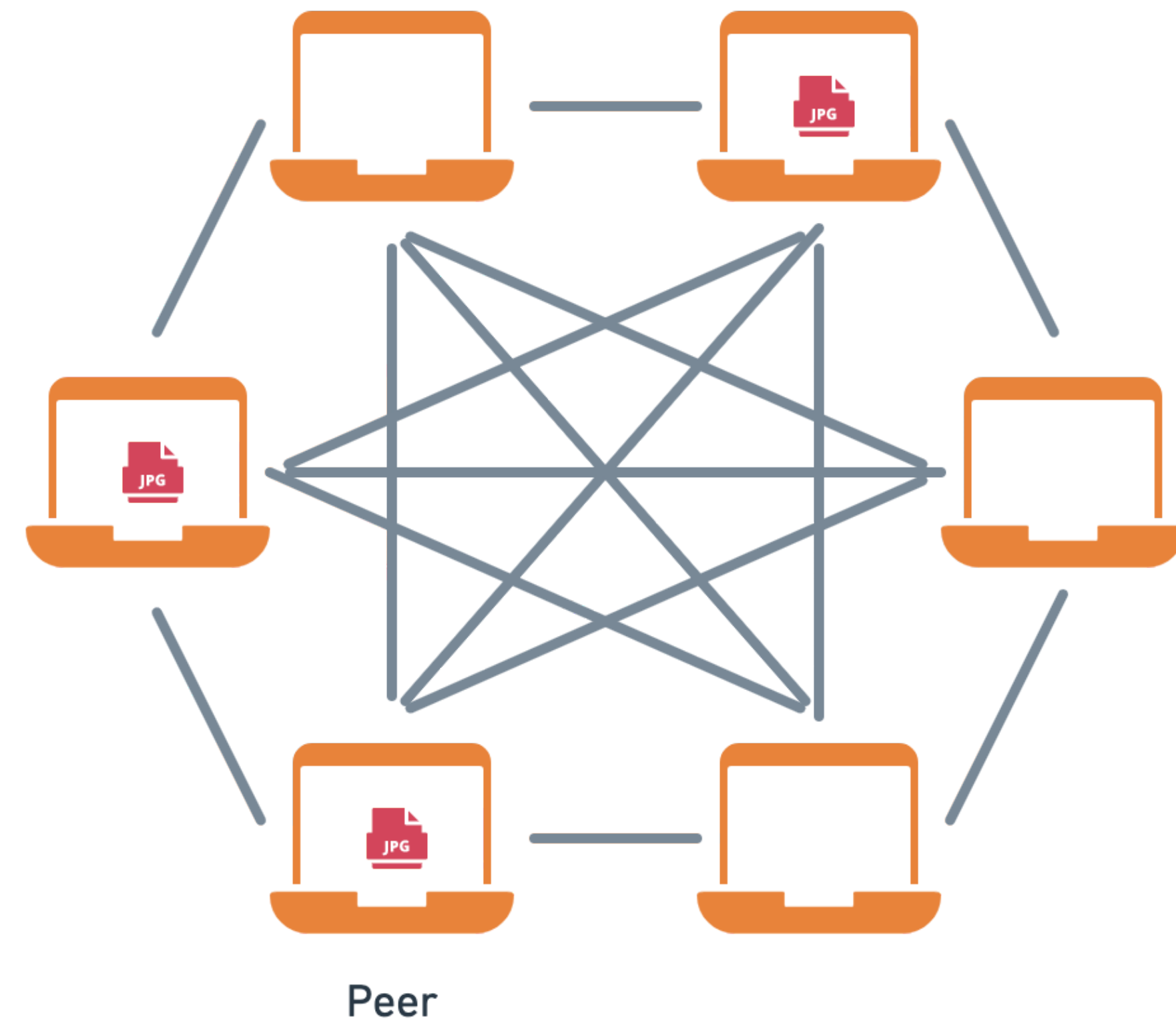
4. Filecoin: decentralised storage market power by a blockchain



# Decentralised Storage

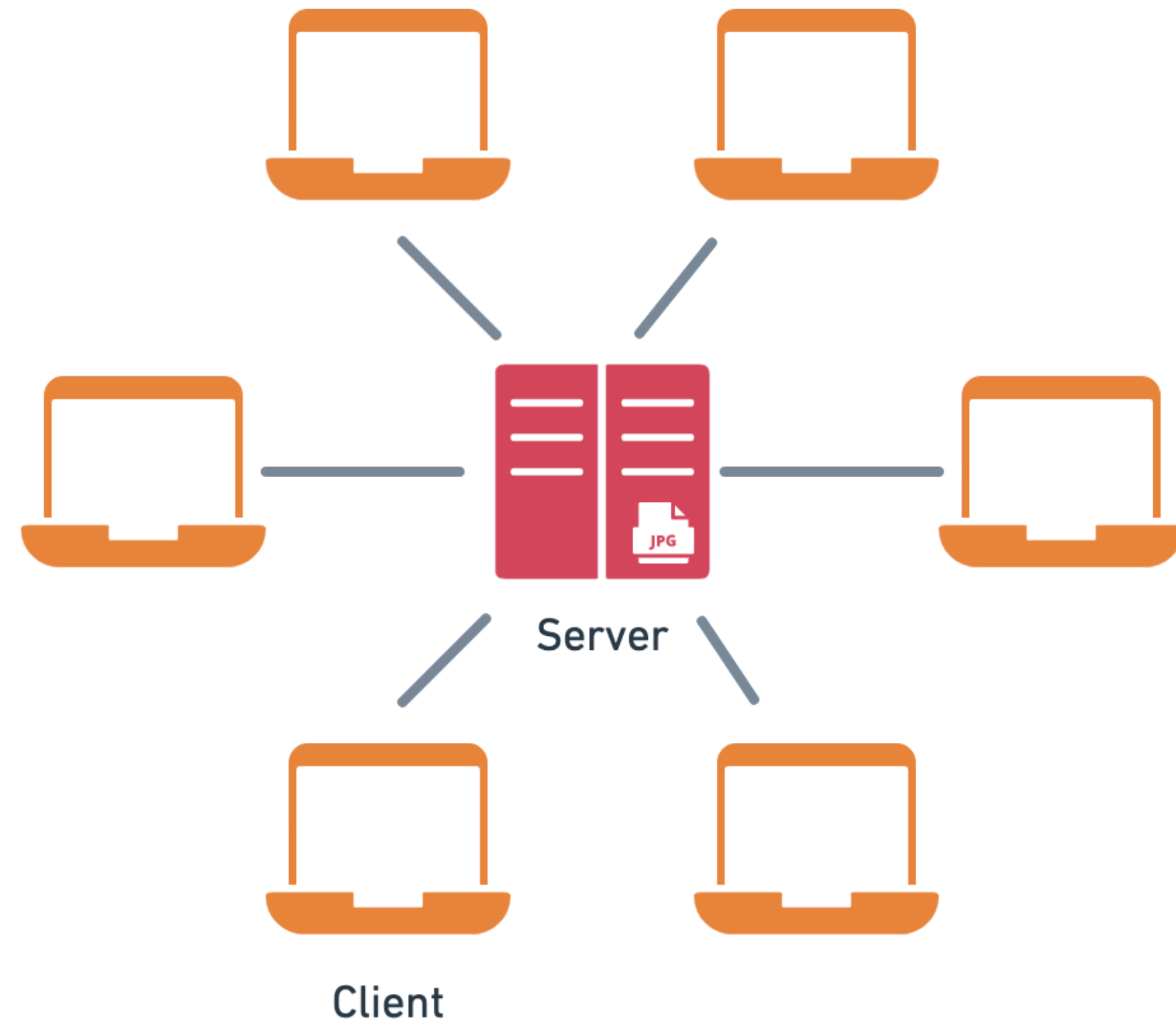


Client-Server model

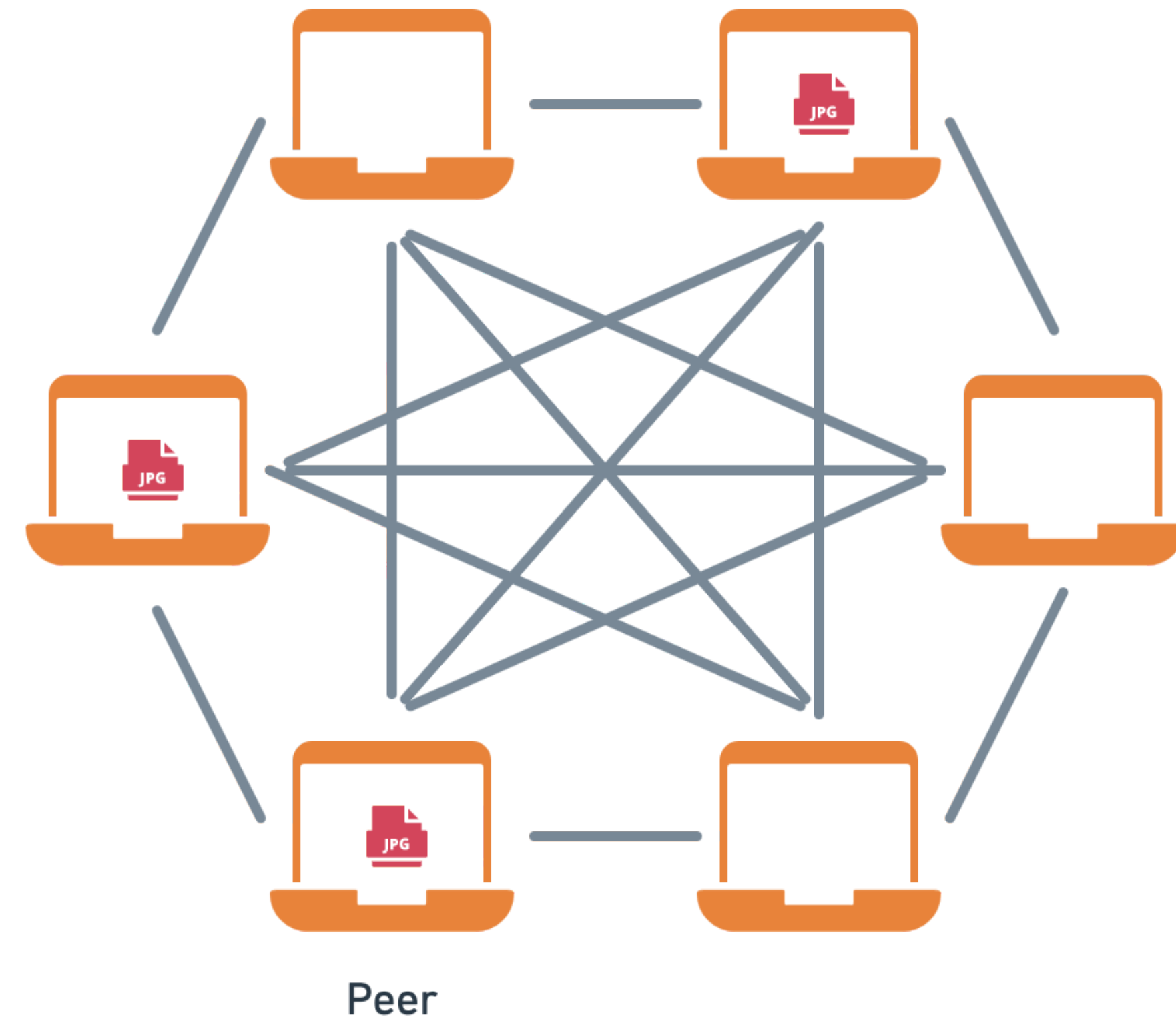


Peer-to-peer model

# Decentralised Storage Market



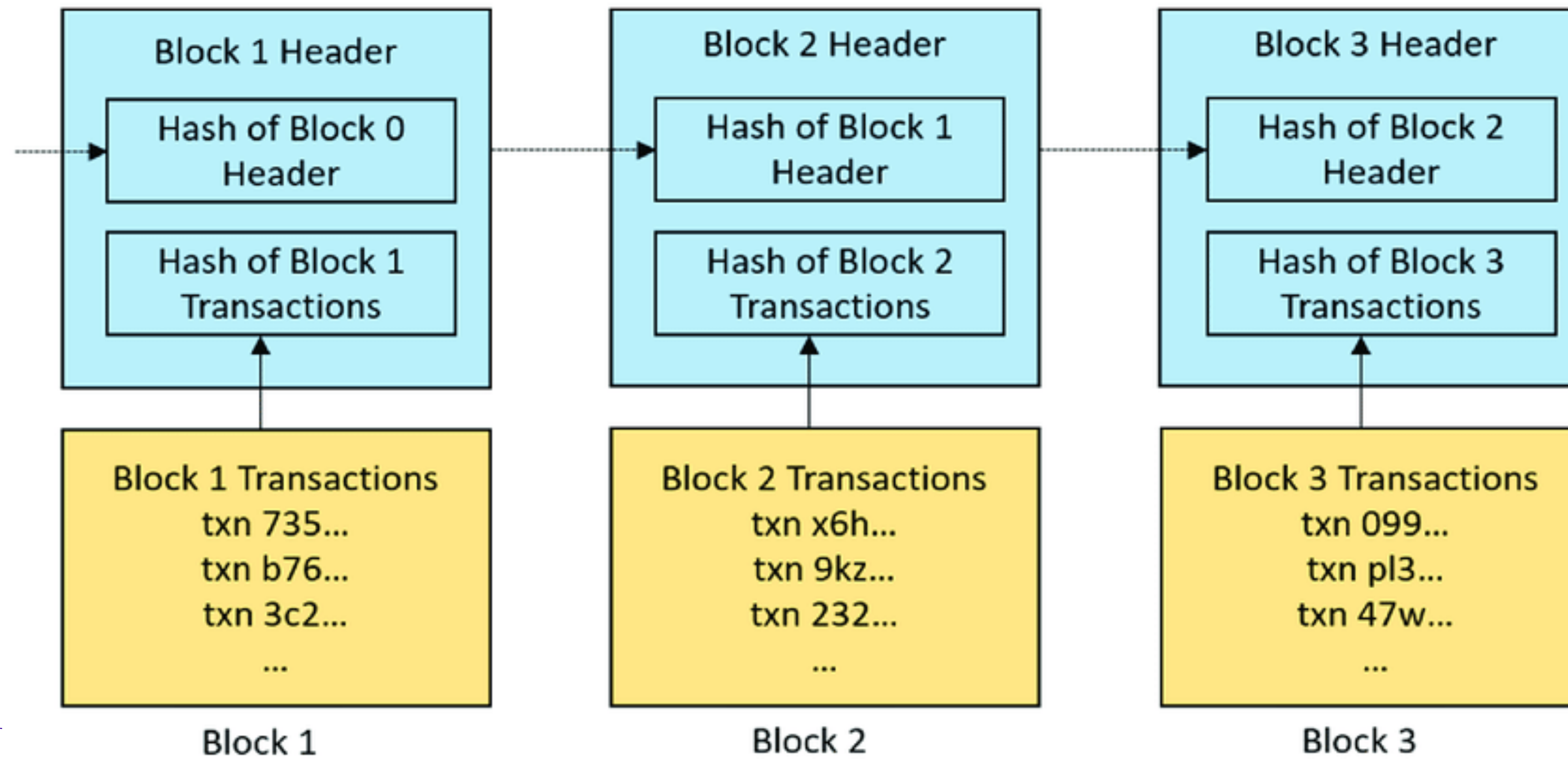
Client-Server model



Peer-to-peer model

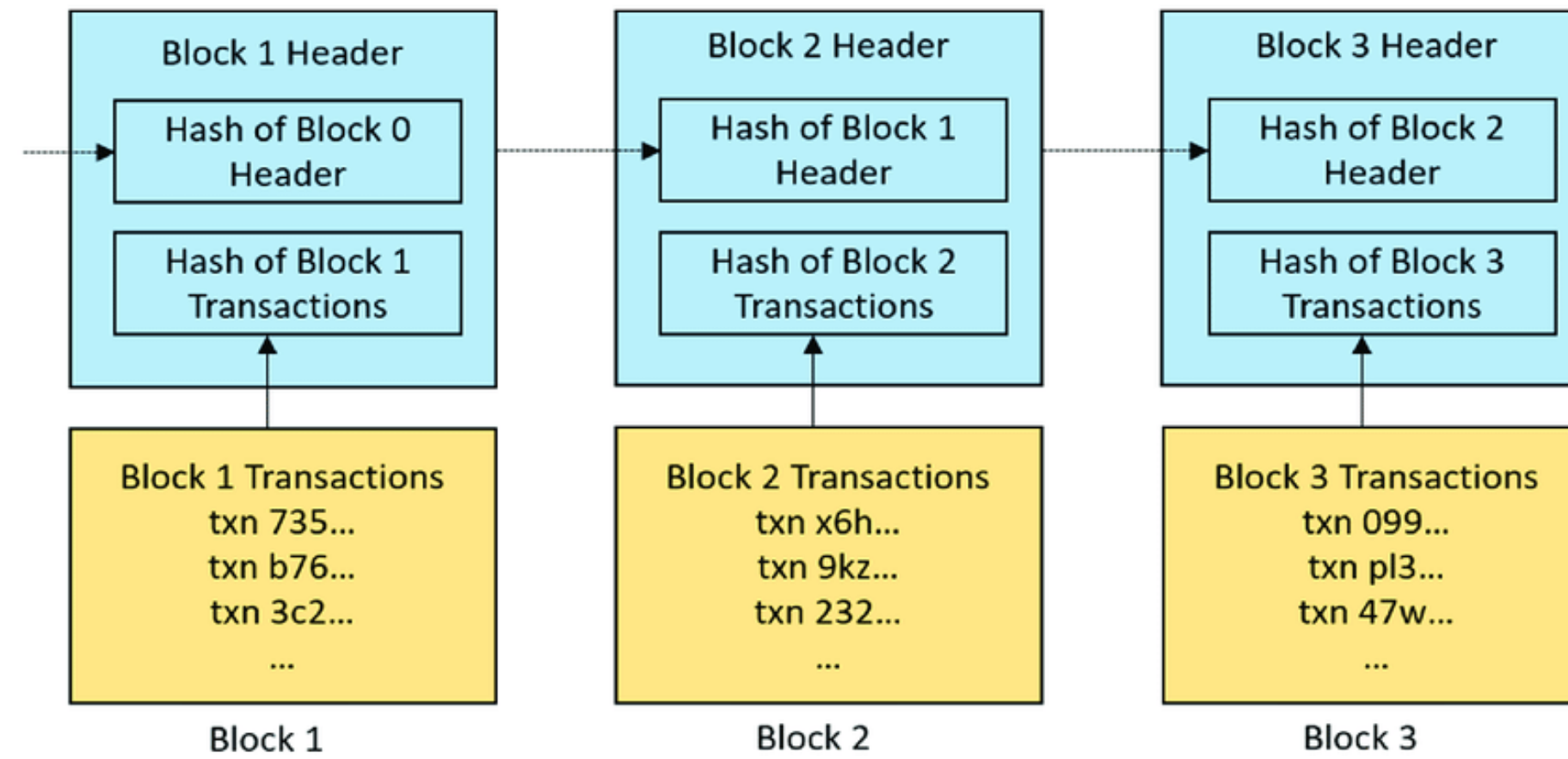
# Blockchain preliminaries

Blockchain = distributed (ie, maintained by a network) ledger, organised in blocks and irreversible



# Blockchain preliminaries

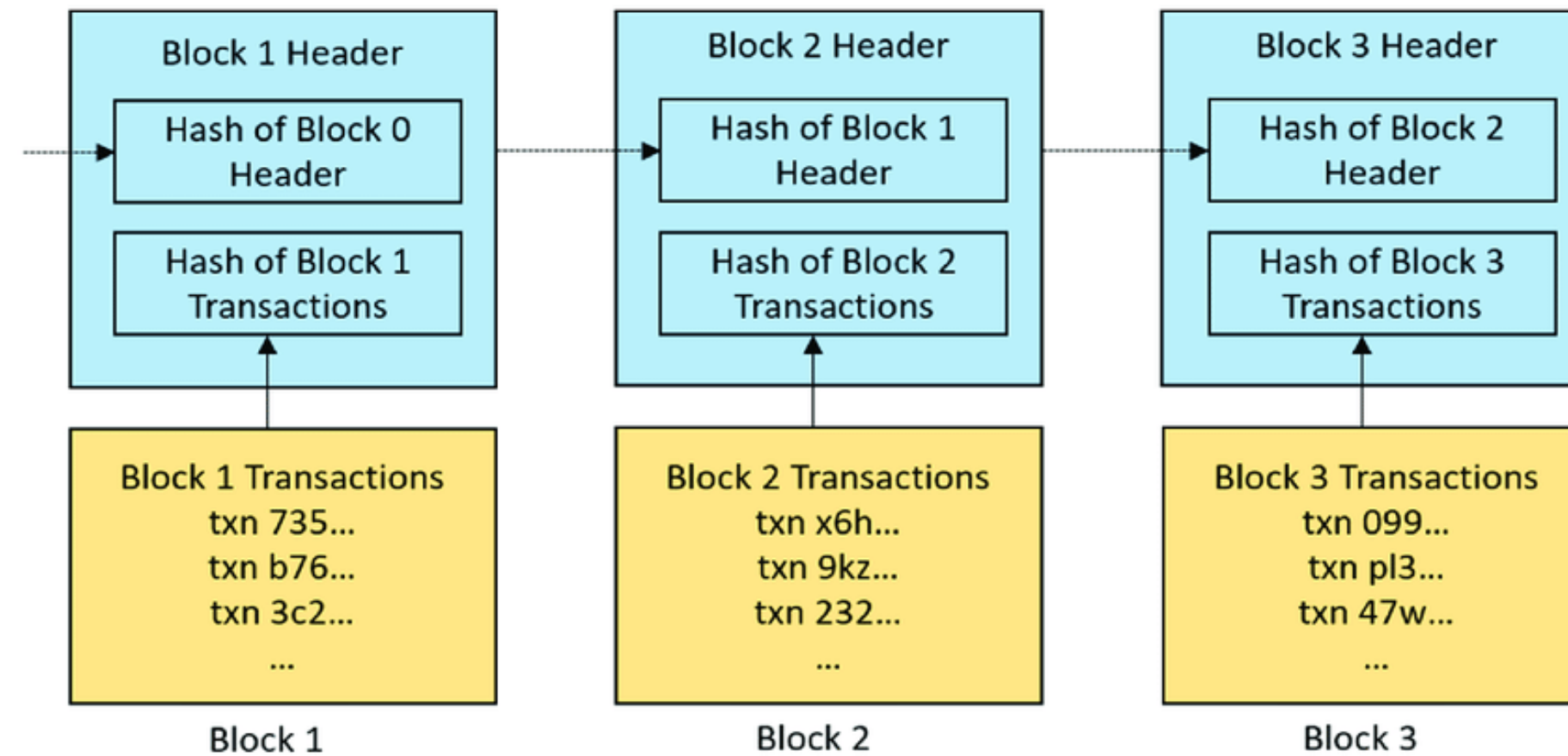
**Critical problem:** How do we agree on the next block?





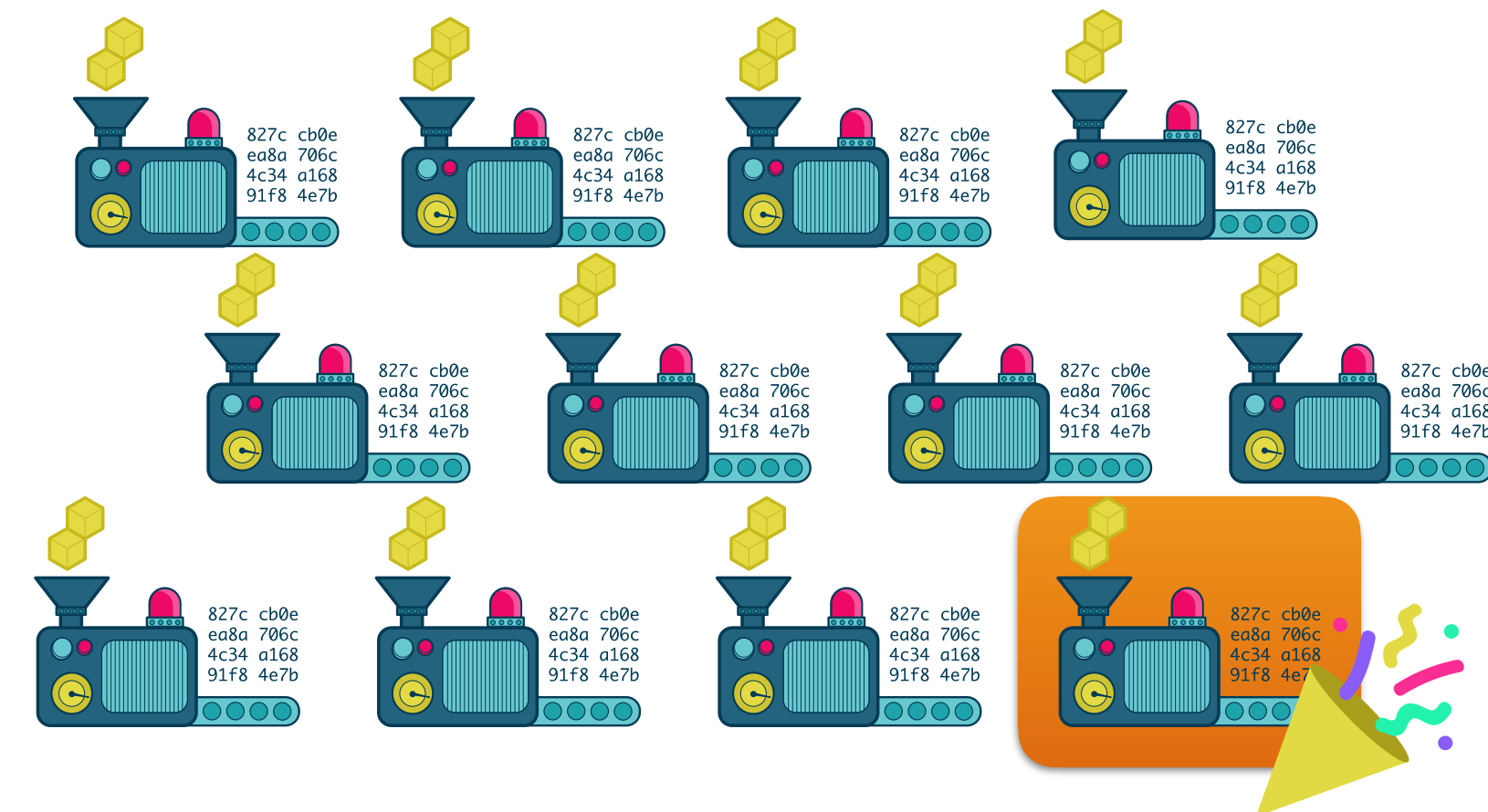
# Blockchain preliminaries

**Critical problem:** How do we agree on the next block?

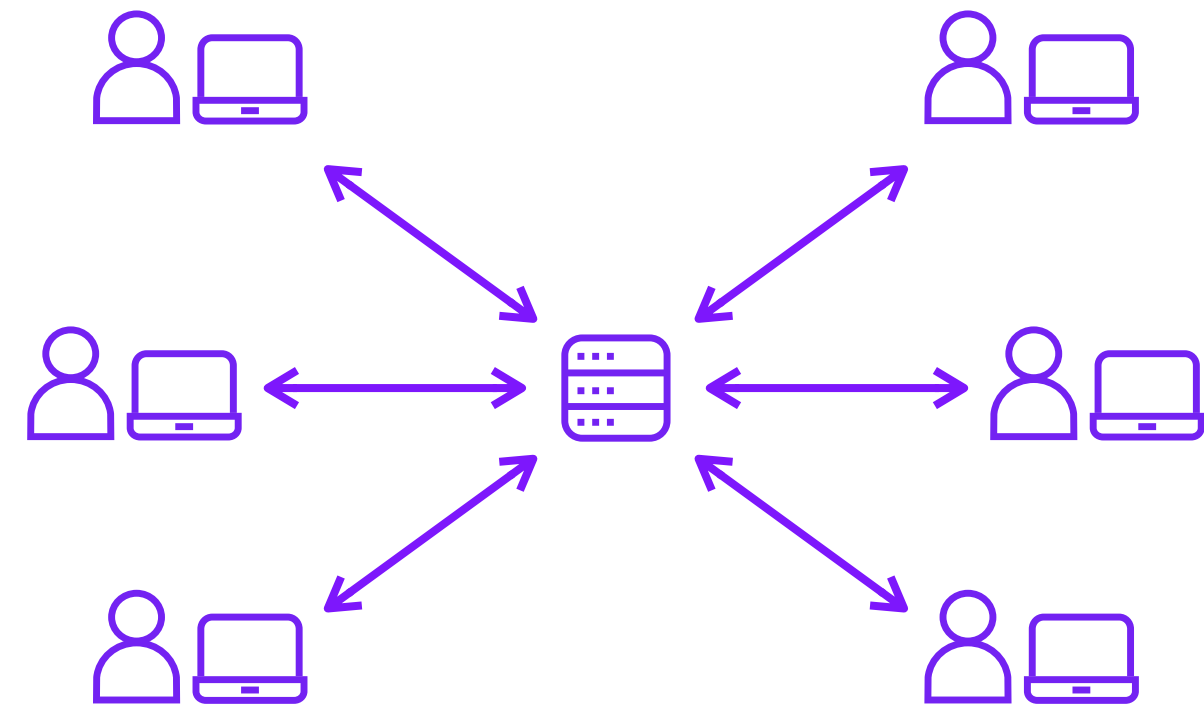


A **lottery** based on computational power chooses one node who creates the next block!

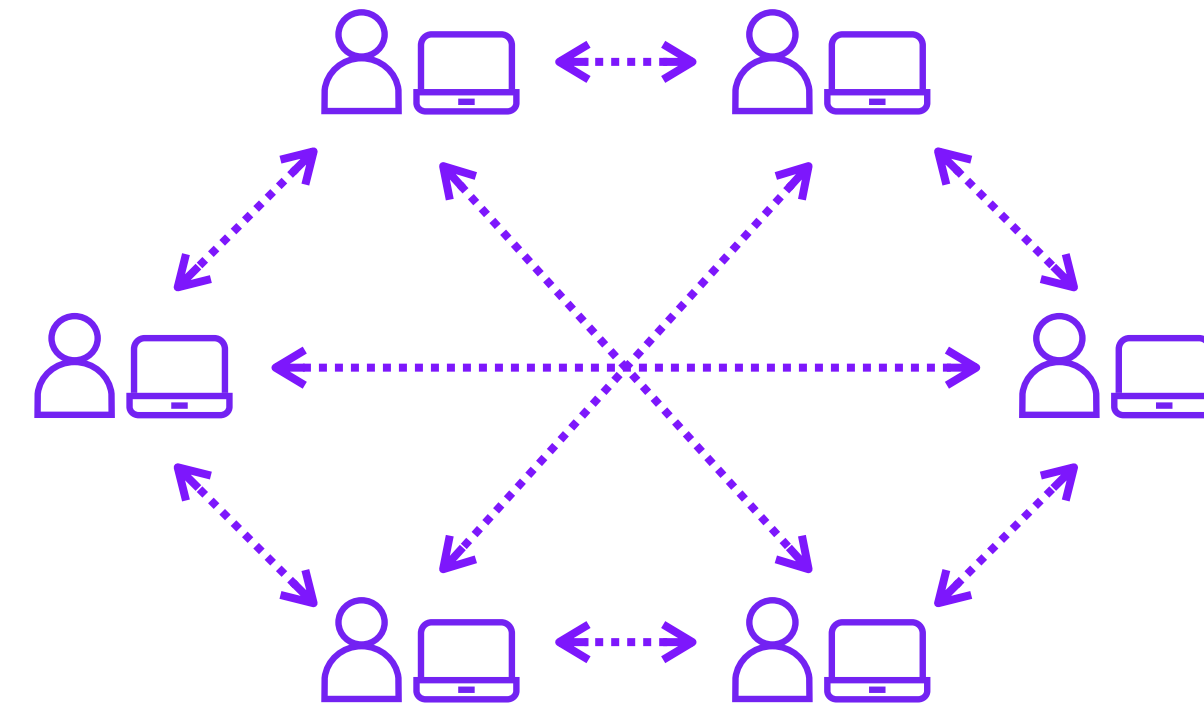
**(Nakamoto Consensus Protocol)**



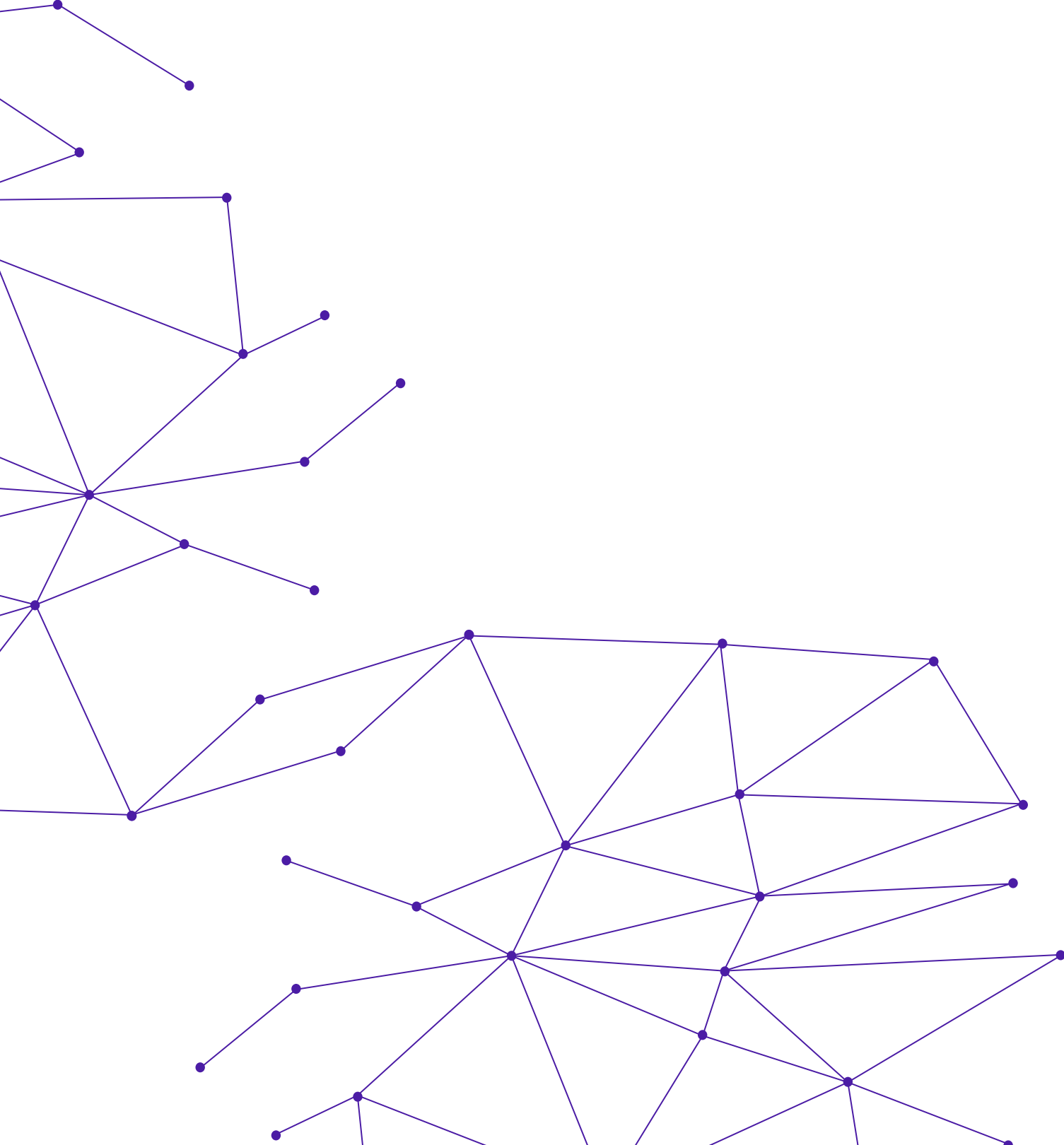
# Can a blockchain give decentralised storage market?



VS






Filecoin: Decentralise the Storage Market



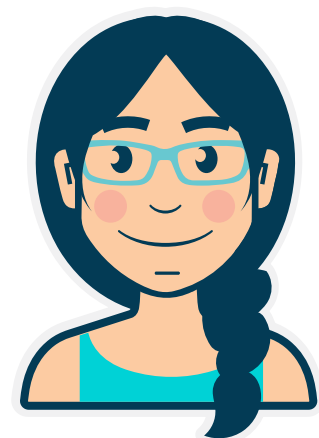
The **Blockchain** replaces the **Trusted Third Party** by keeping track of everything going on in the network

# Blockchains and Storage

- ▶ Data availability: Store transactions data (from Ethereum)
  - Example: [celestia.org](https://celestia.org) 
- ▶ Data in the blockchain's blocks
  - Example:  [arweave.org](https://arweave.org)
- ▶ Blockchain and smart contracts are used for deals and payments
  - Example: [sia.tech](https://sia.tech)
- ▶ Data Storage is the underlying resource for the blockchain consensus
  - Example: [filecoin](https://filecoin.io) 

# How PoS is different from PoW?

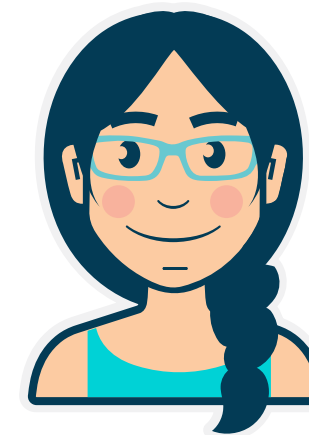
## Proof of Space (PoS)



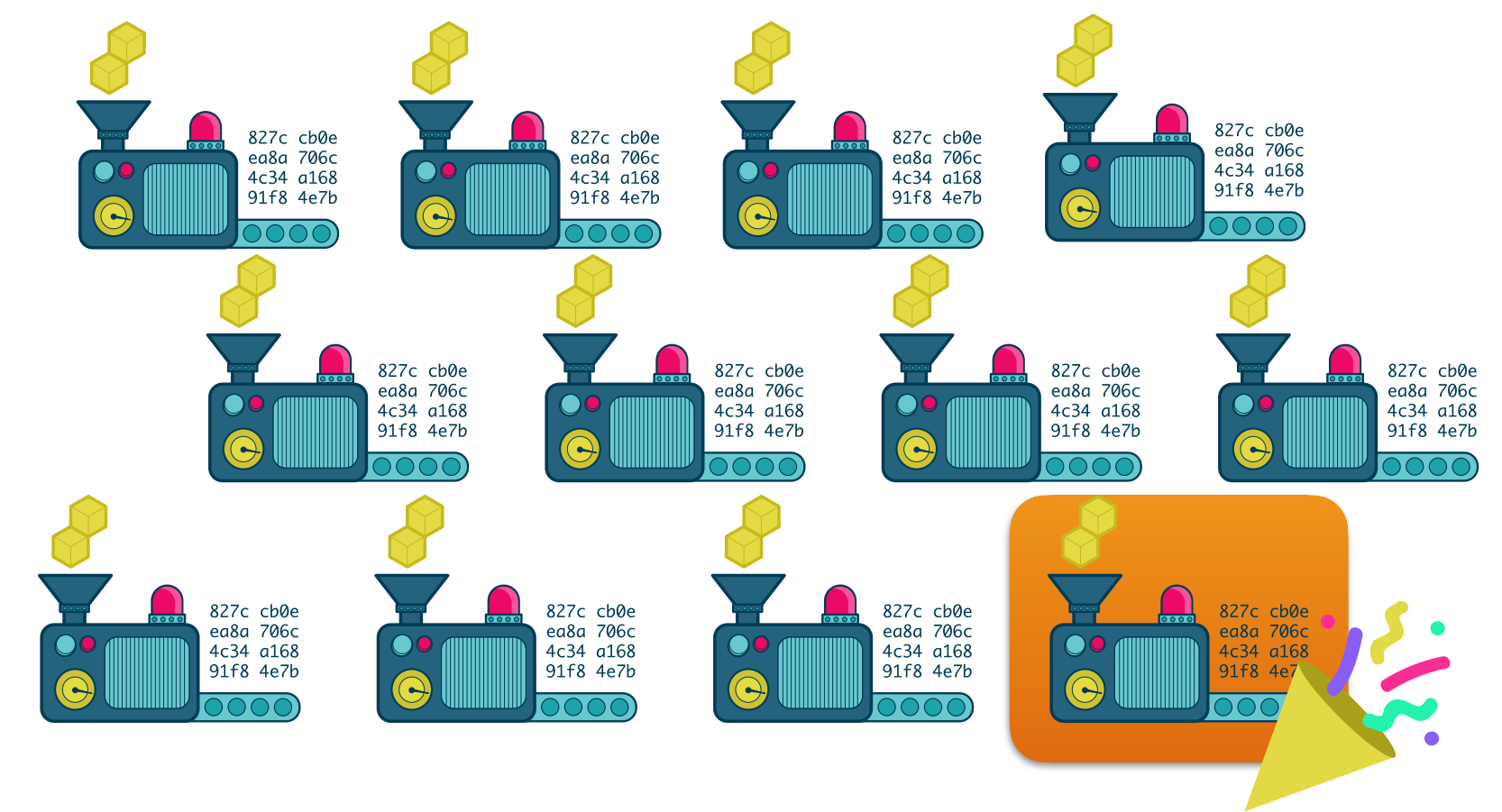
Alice



## Proof of Work (PoW)



Alice



In **PoS** the underlying resource is (persistent) storage rather than computation.

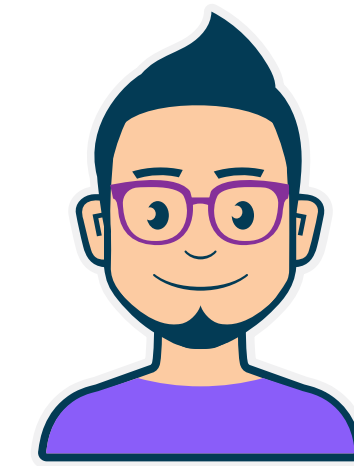
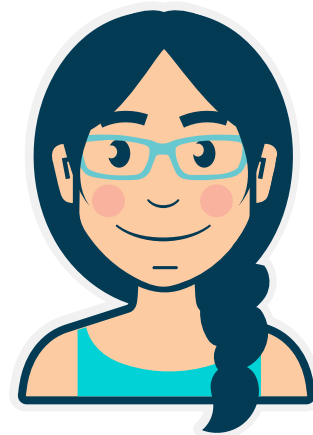
*A more efficient and green alternative to PoW in Nakamoto-Style*

*Consensus Protocols!*

# What is a Proof of (Useful) Space?

## Proof of Space (PoS)

Alice  
“I am storing data  
D of size N!”



Bob  
“True/False!”

Security property:  
*Bob can check that Alice persistently stores  
the (incompressible) data D of size N!*

# What is a Proof of Space?

## Proof of Space (PoS)

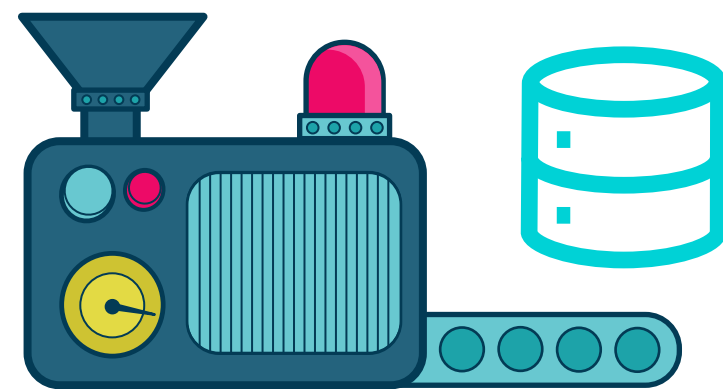
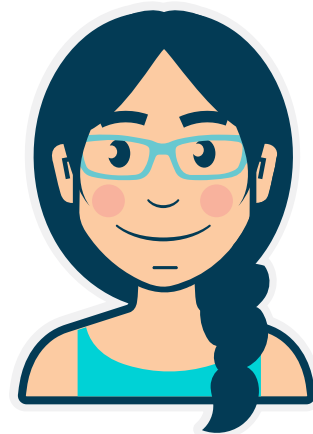
### Two phases protocol:

1. **Initialization** (one-time setup)
2. **Execution** (repeated audit phase)

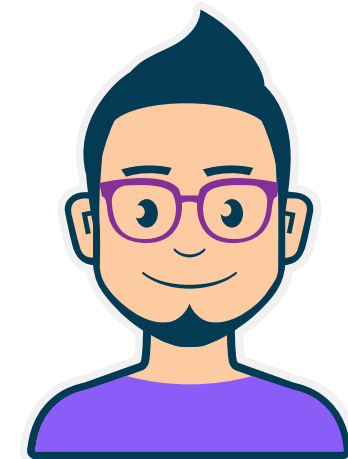


**Initialization:**

Alice



**Advice,  
Enc(data)**



Bob

# What is a Proof of Space?

## Proof of Space (PoS)

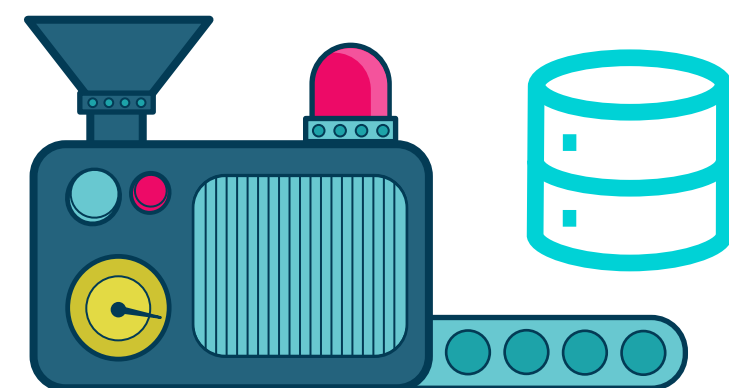
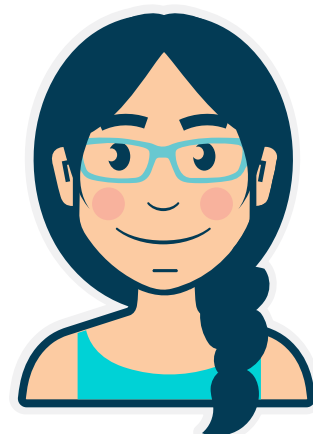
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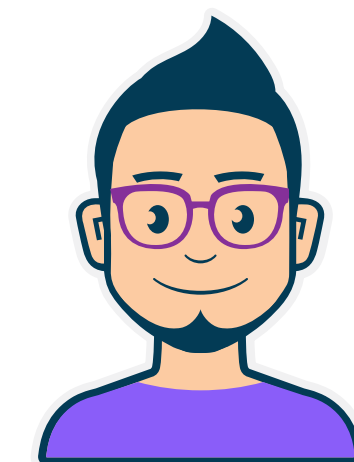


Initialization:

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Advice,  
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Bob

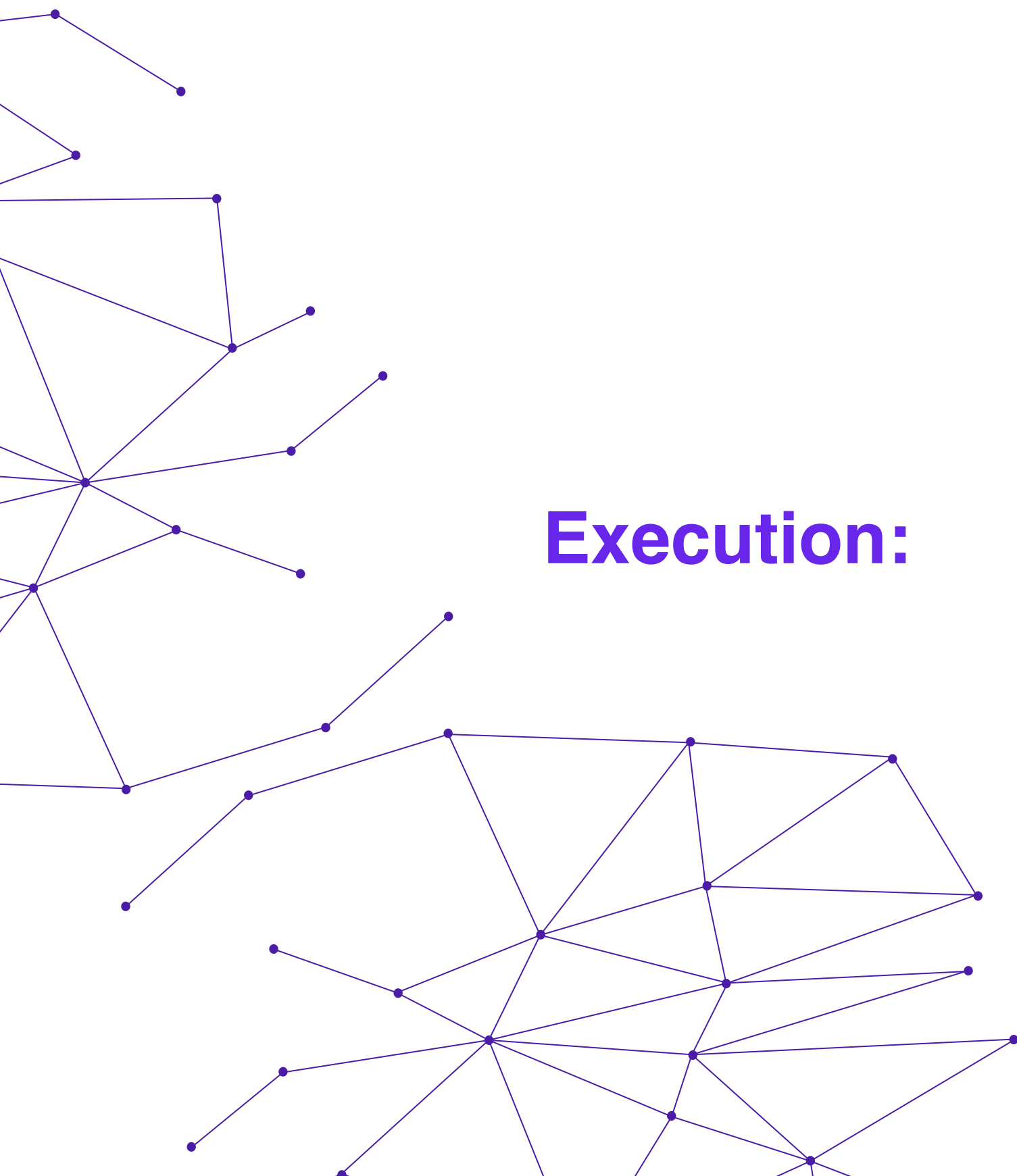
Tag(  )

# What is a Proof of Space?

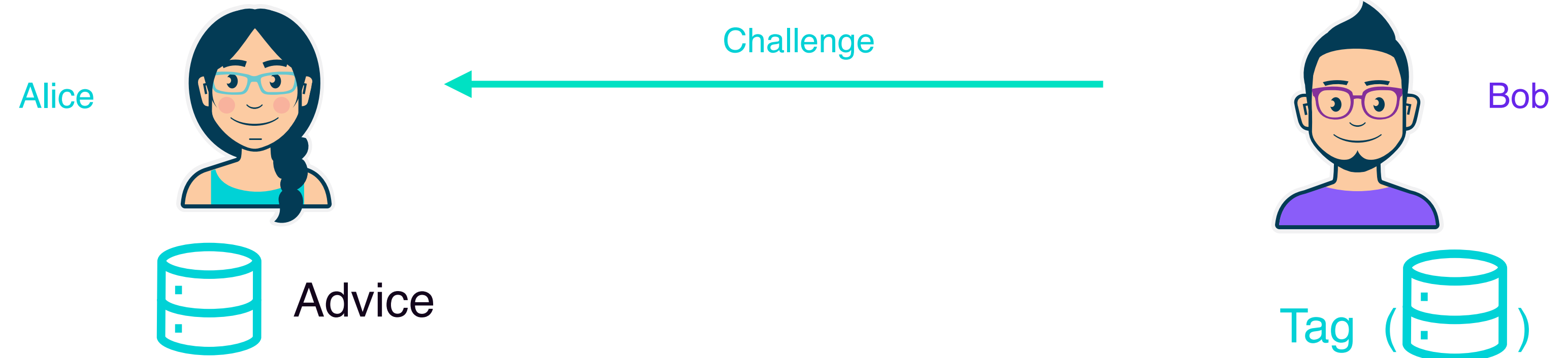
## Proof of Space (PoS)

### Two phases protocol:

1. **Initialization** (one-time setup)
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### Execution:





# What is a Proof of Space?

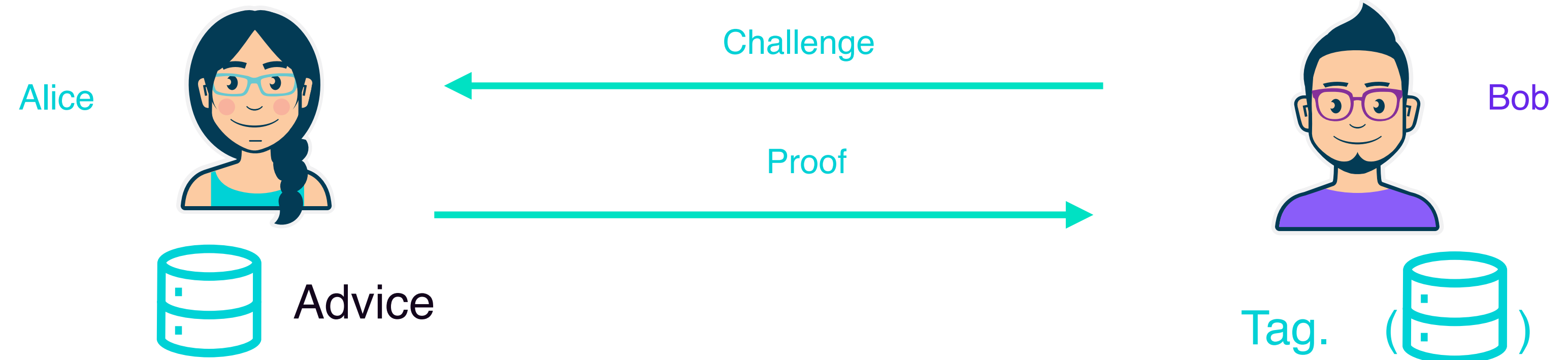
## Proof of Space (PoS)

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### Execution:

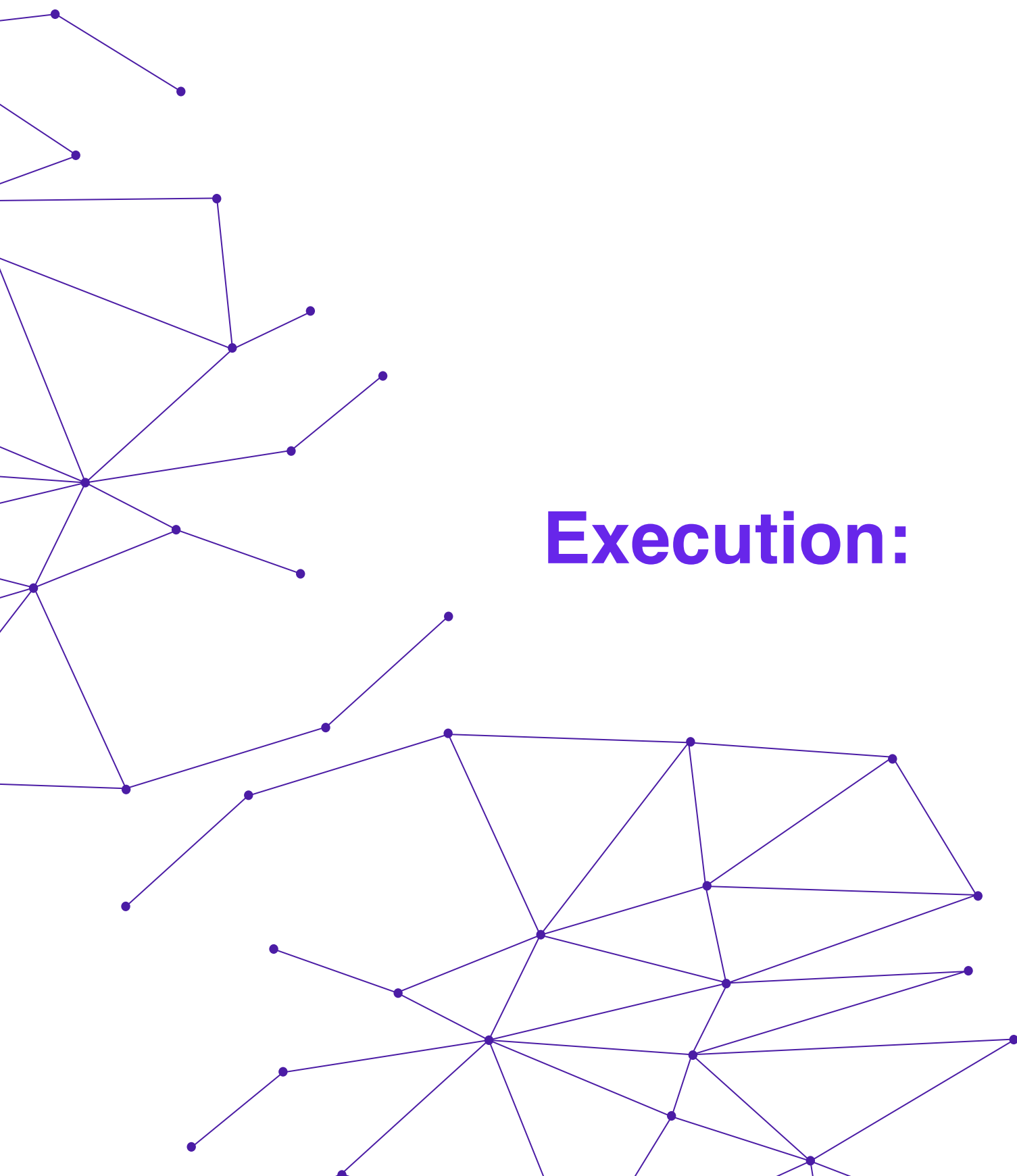


# What is a Proof of Space?

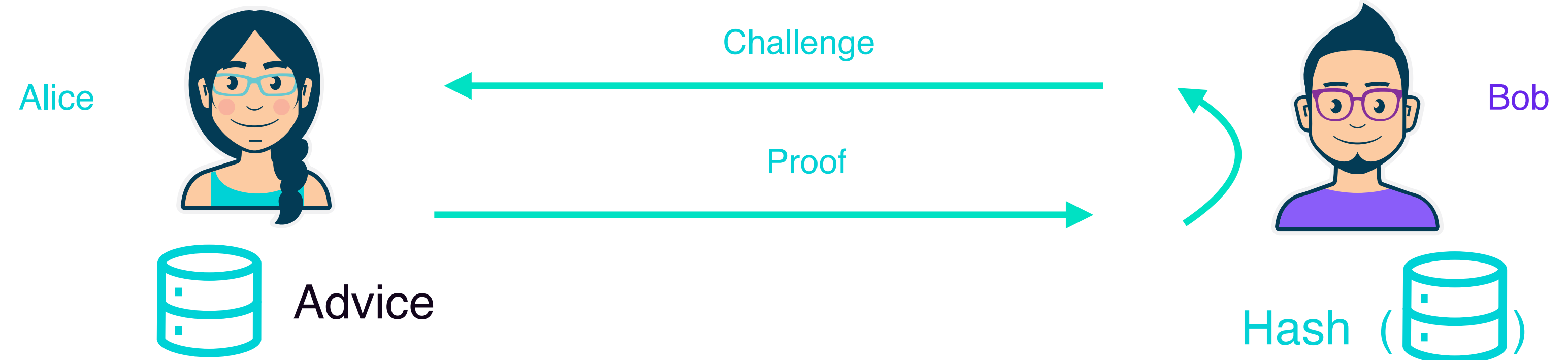
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### Execution:



# What is a Proof of Space?

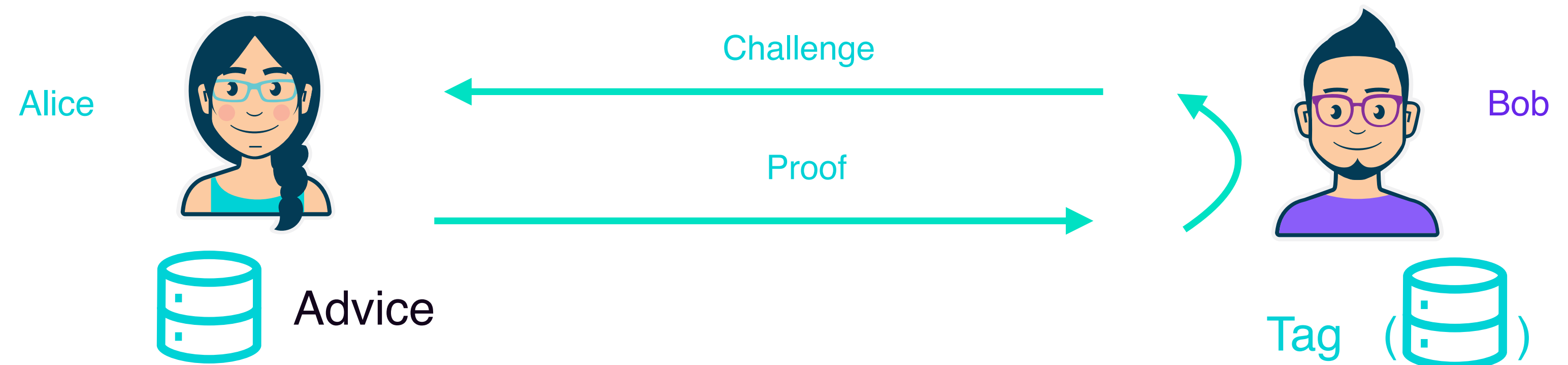
## Proof of Space (PoS)

### Two phases protocol:

1. **Initialization** (one-time setup)

2. **Execution** (repeated audit phase)

### Execution:



### Soundness (informal, fix eps)

if Alice stores  $N'$  then she must do a computation of  $T(N')$  steps in order to produce a convincing proof (for any  $N' < \text{eps} * N$ ).

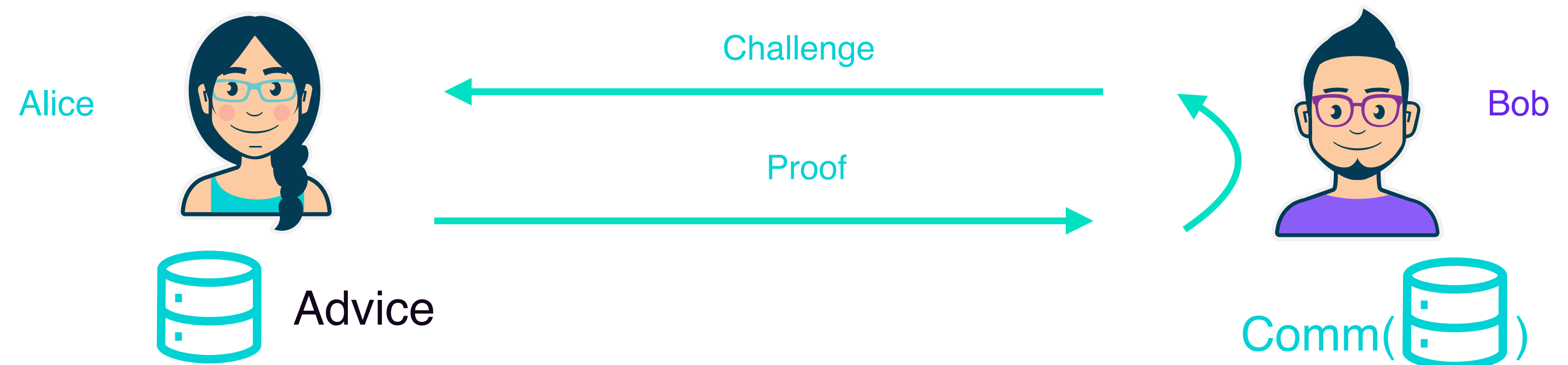
# Question 1: What is a Proof of Space?

## Proof of Space (PoS)

### Two phases protocol:

1. **Initialization** (one-time setup)
2. **Execution** (repeated audit phase)

### Execution:



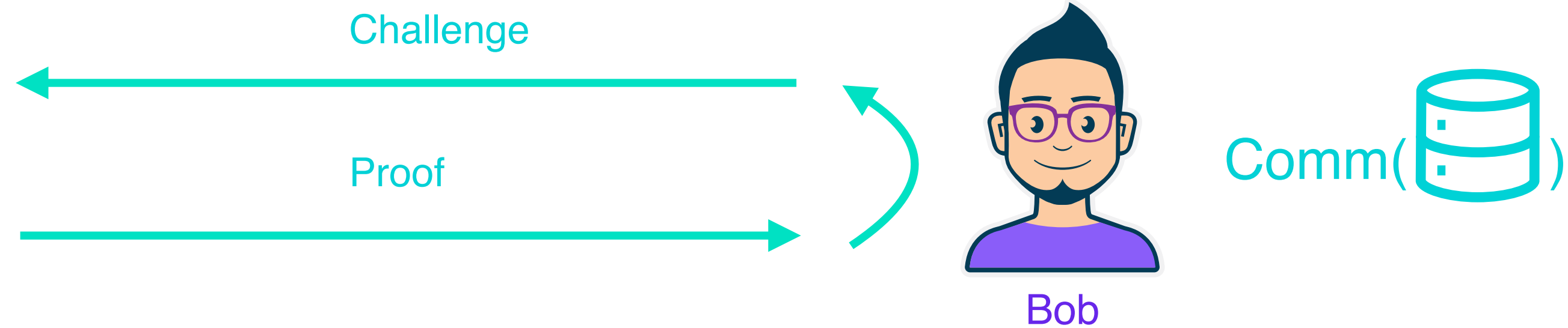
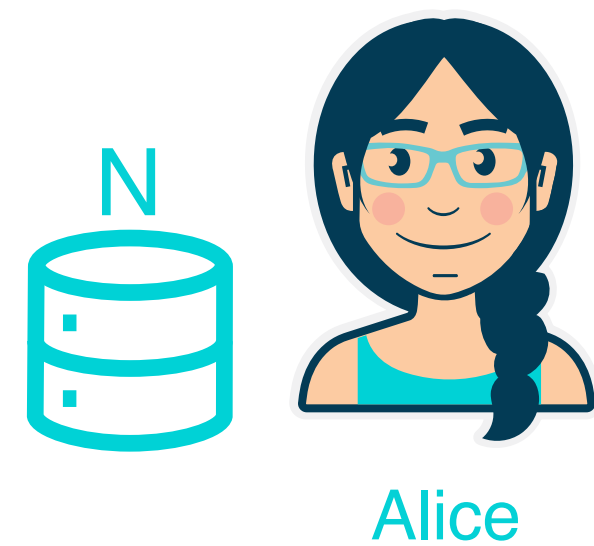
The "*T-step computation*" is infeasible/irrational  
=> Alice is persistently storing the advice

# What is a Proof of Space?



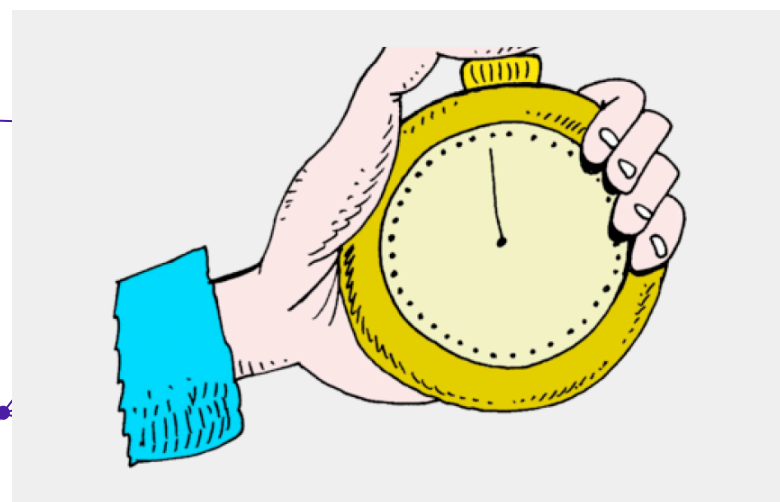
## Proof of Space (PoS)

Execution



What does “T-step computation is infeasible” mean in practice?

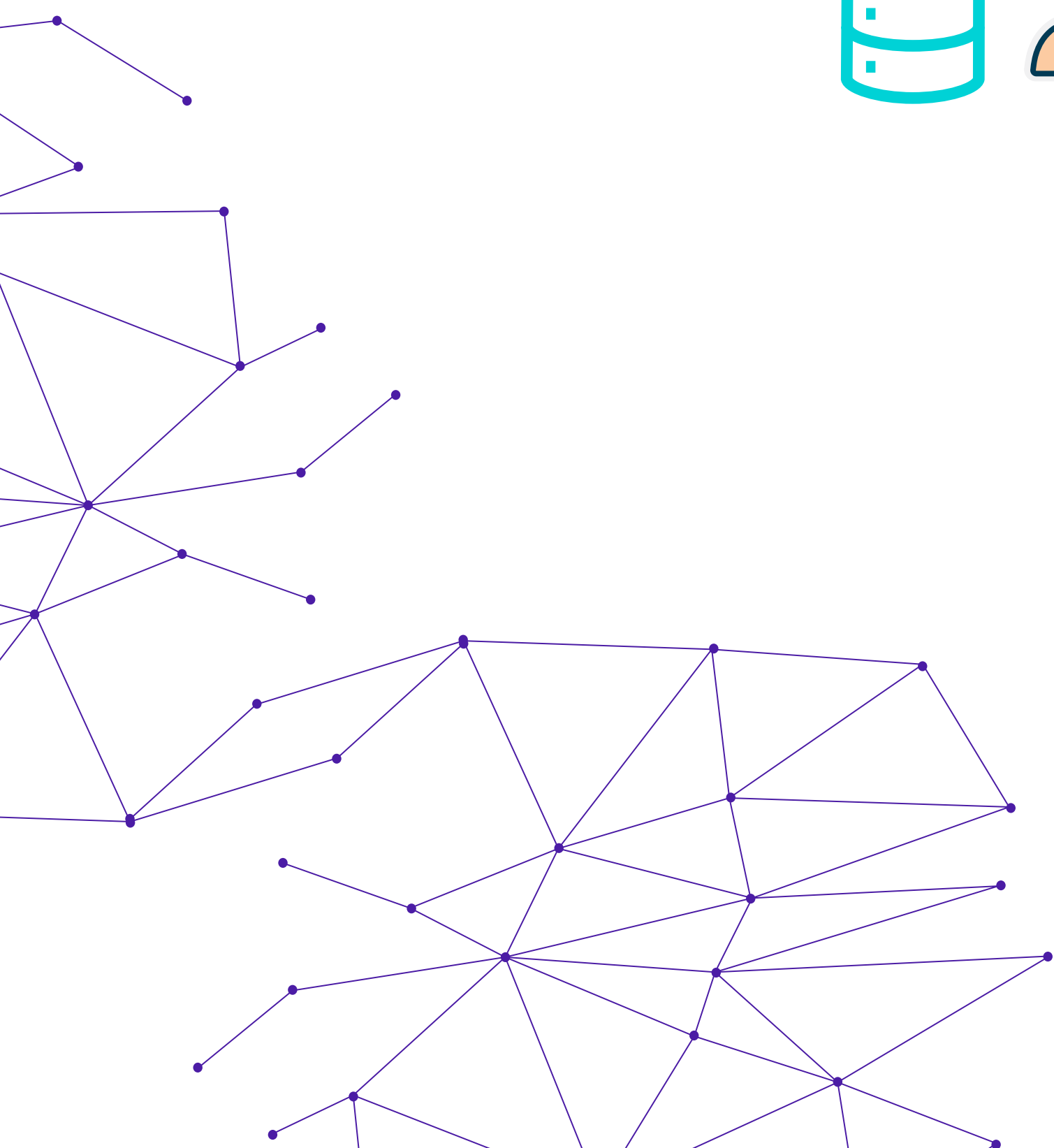
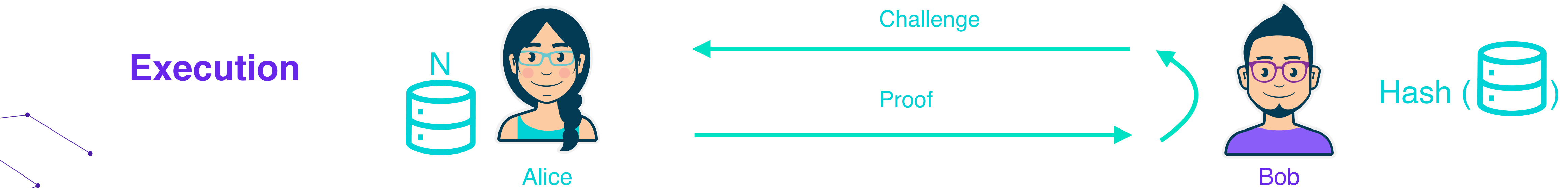
**Latency model:** Alice must answer challenges in a limited window of time for the proof to be valid. We need a **timing heuristic assumption** to translate T-steps in t seconds!



# What is a Proof of Space?



## Proof of Space (PoS)



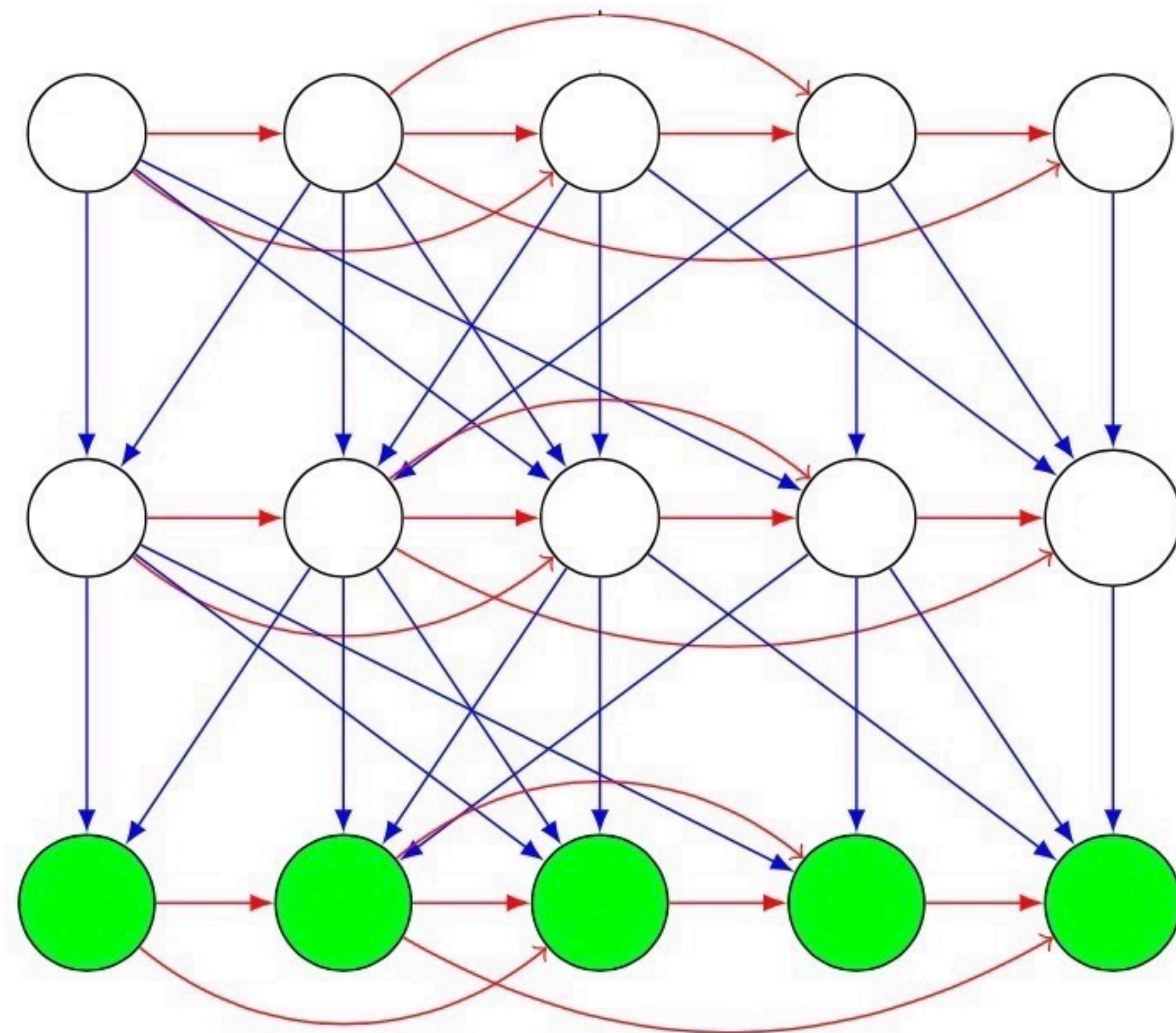
What does “T-step computation is irrational” mean in practice?

**Cost Model:** Alice chooses to store because T steps are more expensive than storing the advice for the time between two consecutive executions (rational prover). We need a **cost heuristic assumption** to translate T-steps in t dollars!



# Graph-labelling based PoS

## Stacked-DRGs graph (aka “SDR”)



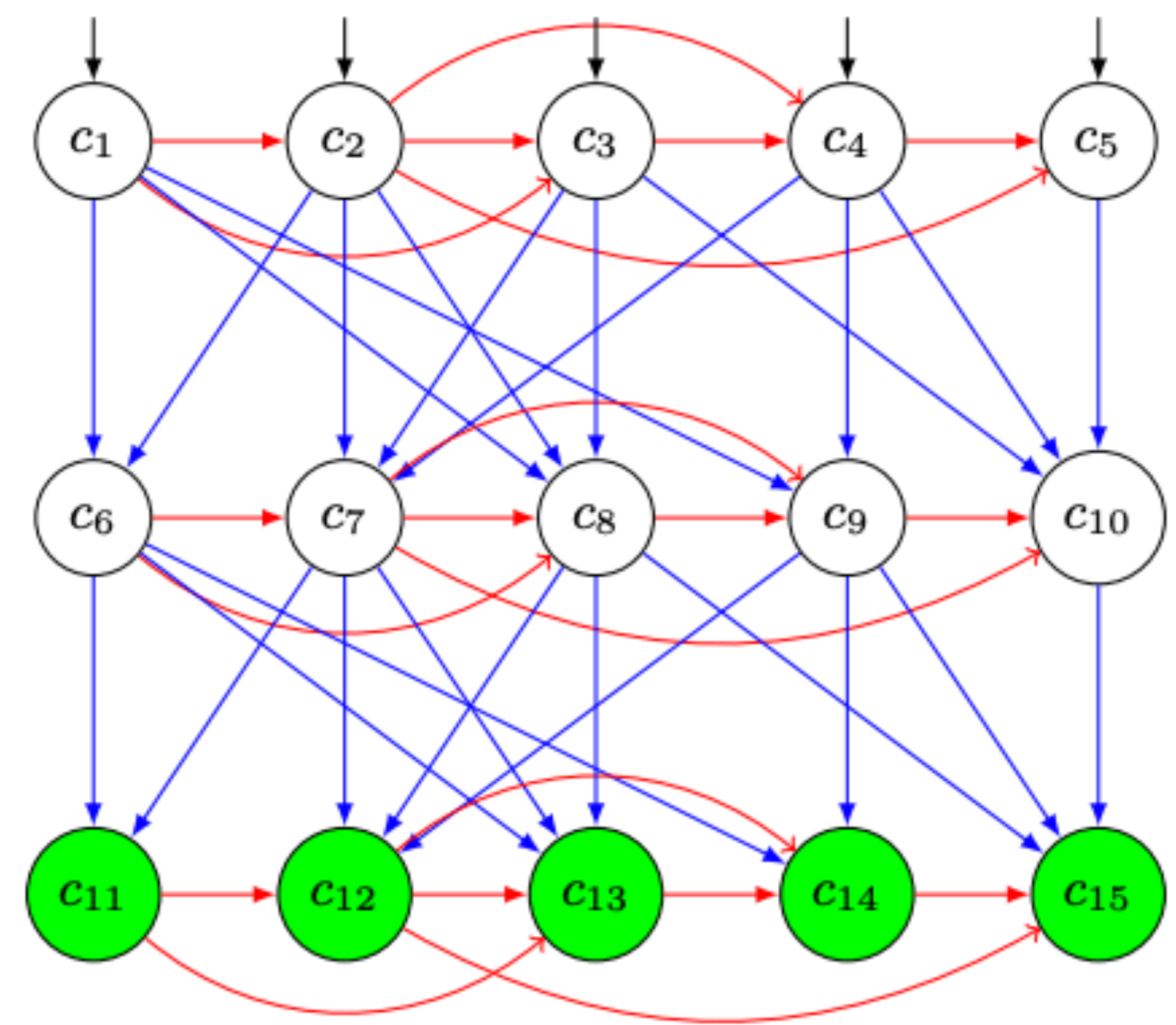
Each layer is a **depth-robust graph** (red edges):  
*each set of 80% of the nodes has a long path (20%)*

Each pair of layers is an **expander** (blue edges):  
*an  $x$  fraction of nodes in the lower layer has a  $2x$  of the nodes of the upper layer as parents ( $x < 1/3$ )*

Ben Fish, Crypto19, “Tight Proofs of Space”

# Graph-labelling based PoS

## Stacked-DRGs graph (aka “SDR”)

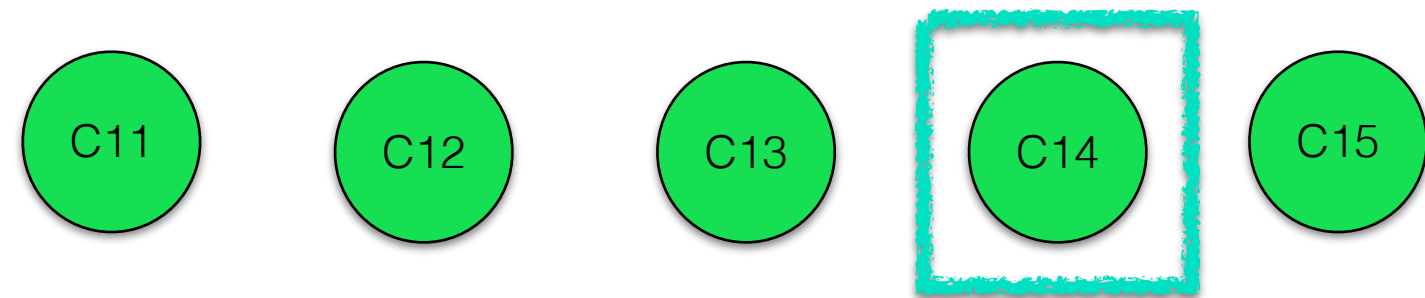


Label of a node:  
 $c_i = \text{Hash}(i \parallel \text{Hash}(D) \parallel c_j \text{ if } j \text{ is a parent of } i)$

$\text{Enc}(D) = \text{labels of green nodes} + D$



# Graph-labelling based PoS

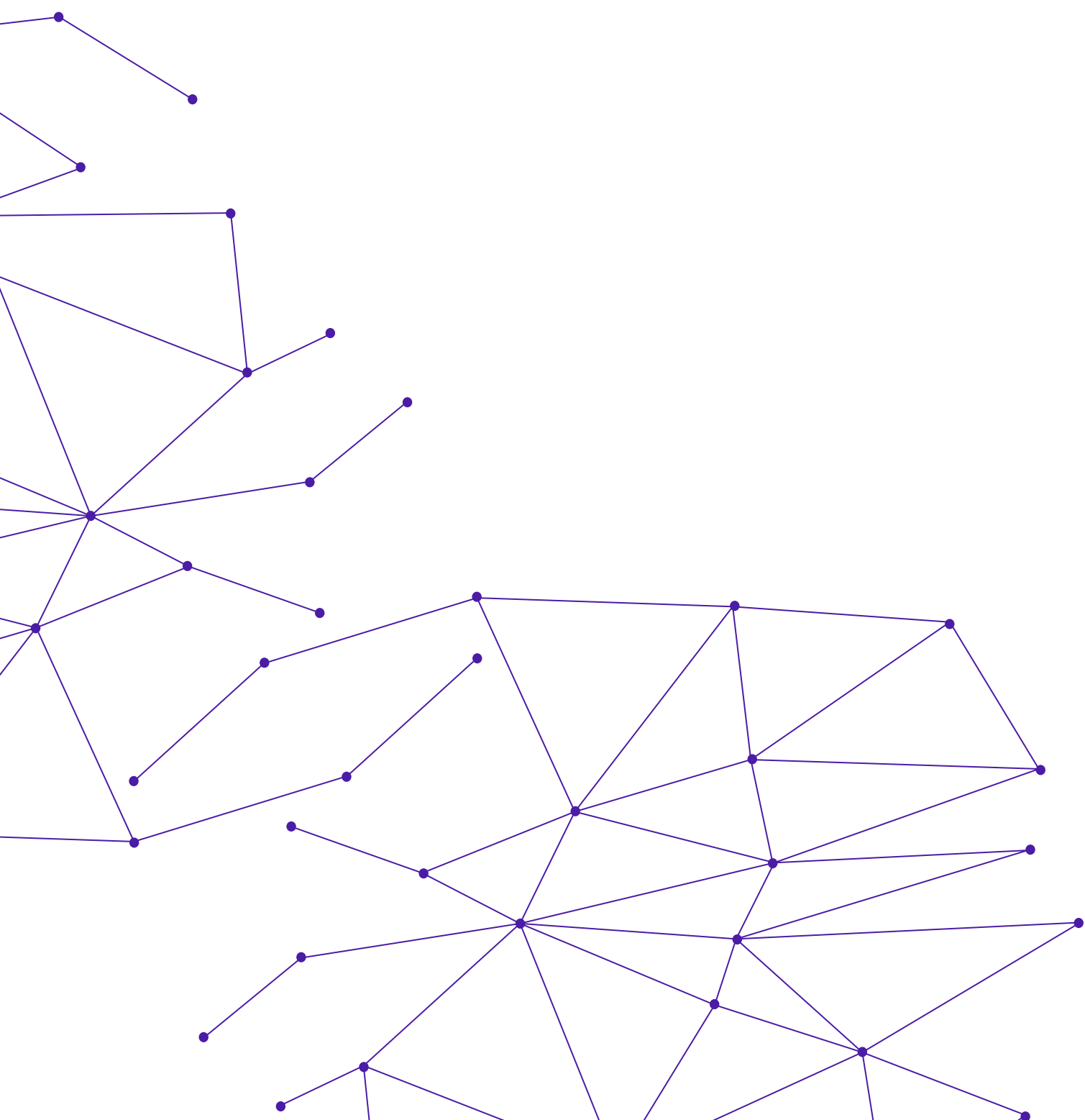


## Execution

Bob periodically chooses a random node

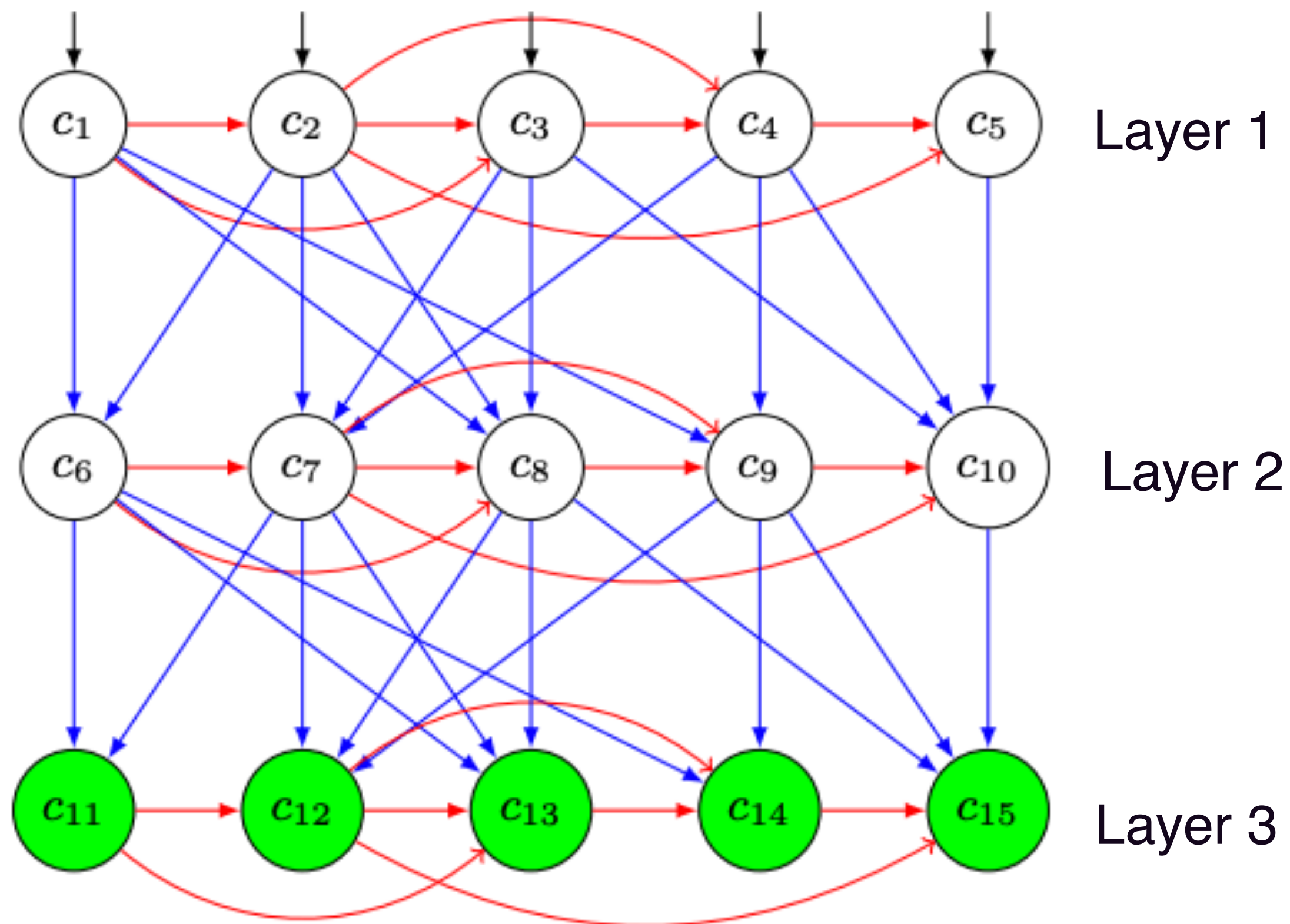
Alice answer with the label (and hash proof)

If she does not store it, must re-compute it!!



# Graph-labelling based PoS

## Stacked-DRGs graph (aka “SDR”)

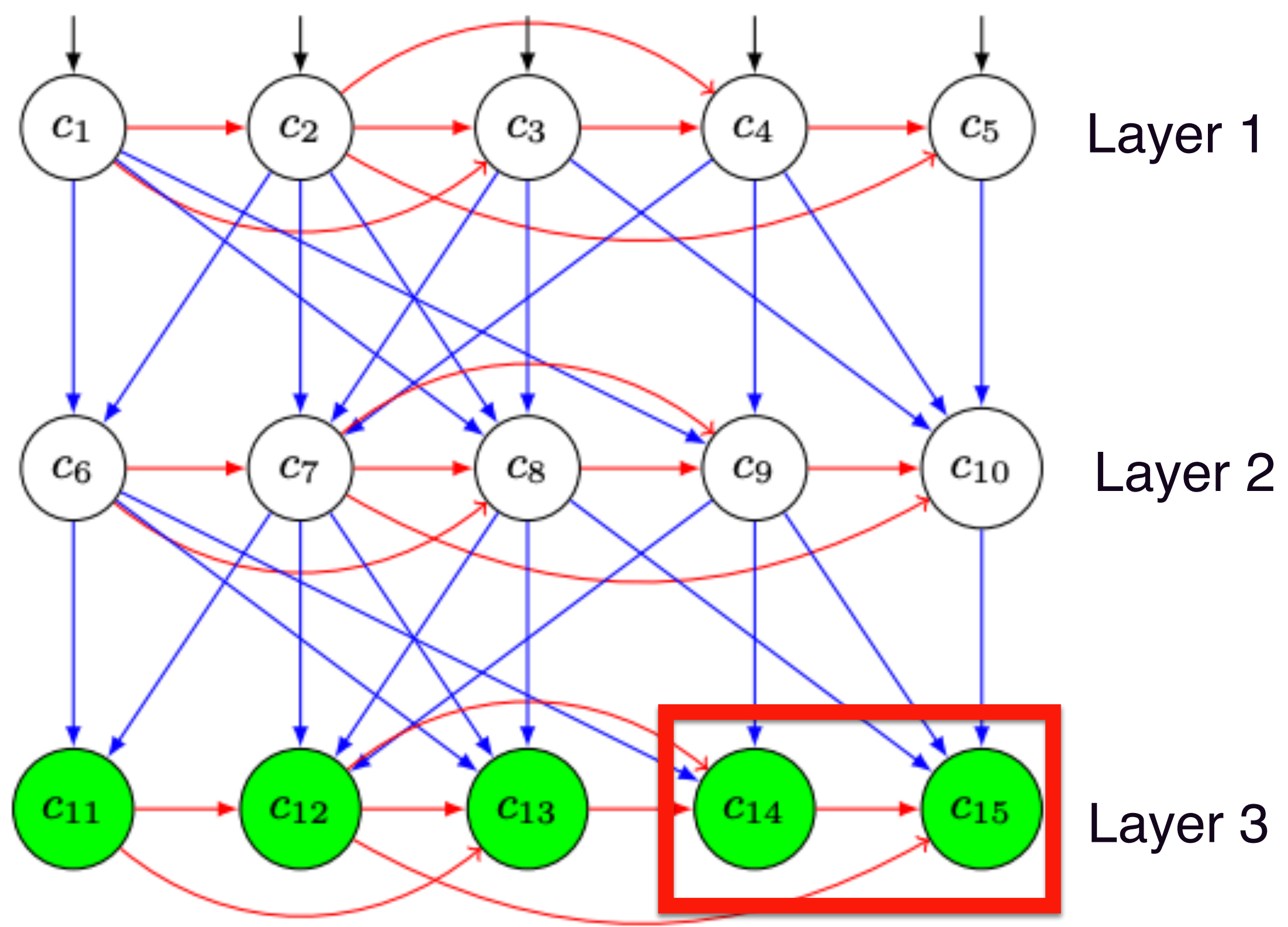


### Security property:

- Re-computing from scratch is slow/expensive because:  
*a random node has “many many” parents (with high prob.)*

# Graph-labelling based PoS

## Stacked-DRGs graph (aka "SDR")

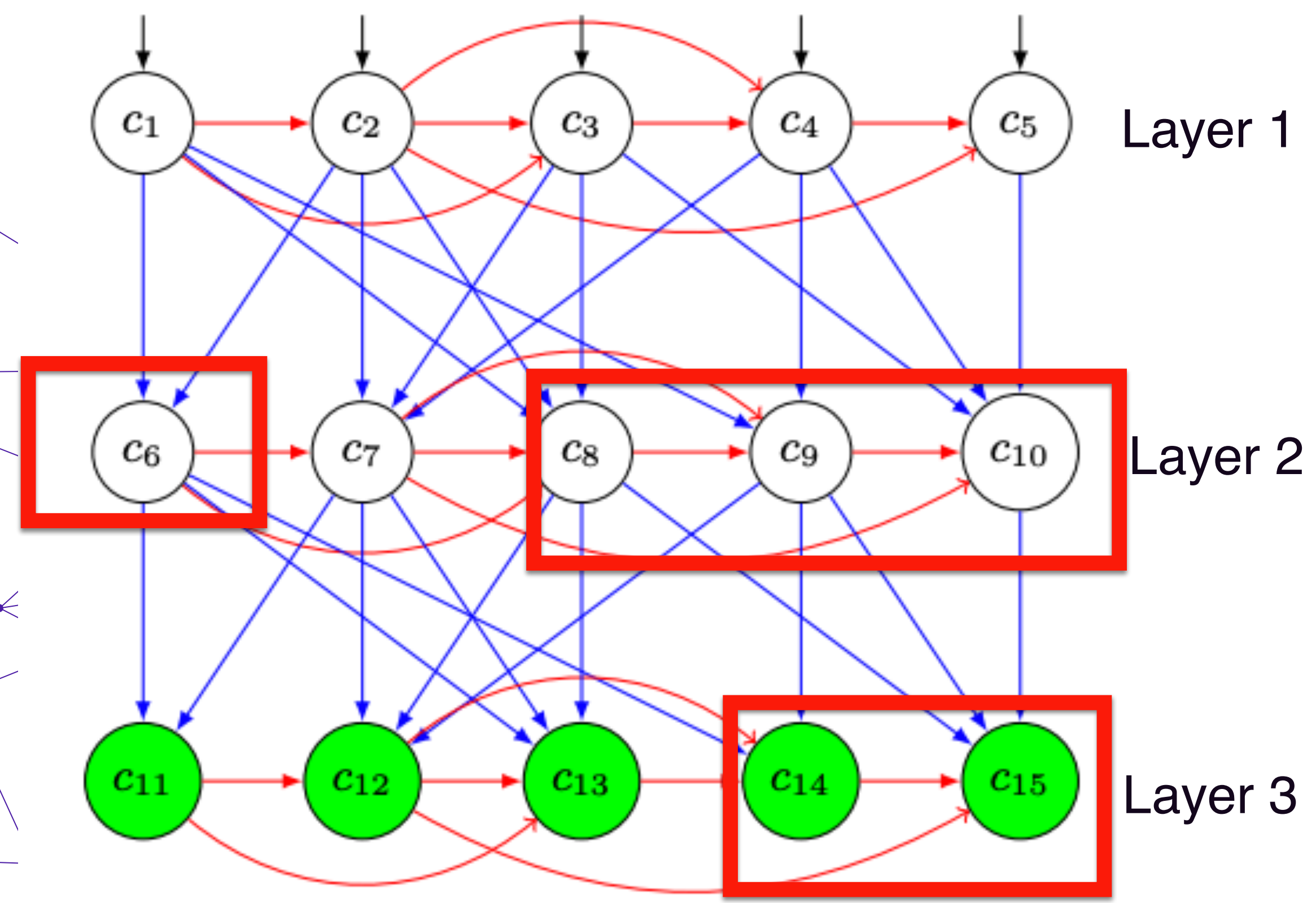


**Malicious Alice:**

Store only  $c_{11}$ ,  $c_{12}$ ,  $c_{13}$   
(Delete  $c_{14}$ ,  $c_{15}$ )

# Graph-labelling based PoS

## Stacked-DRGs graph (aka "SDR")

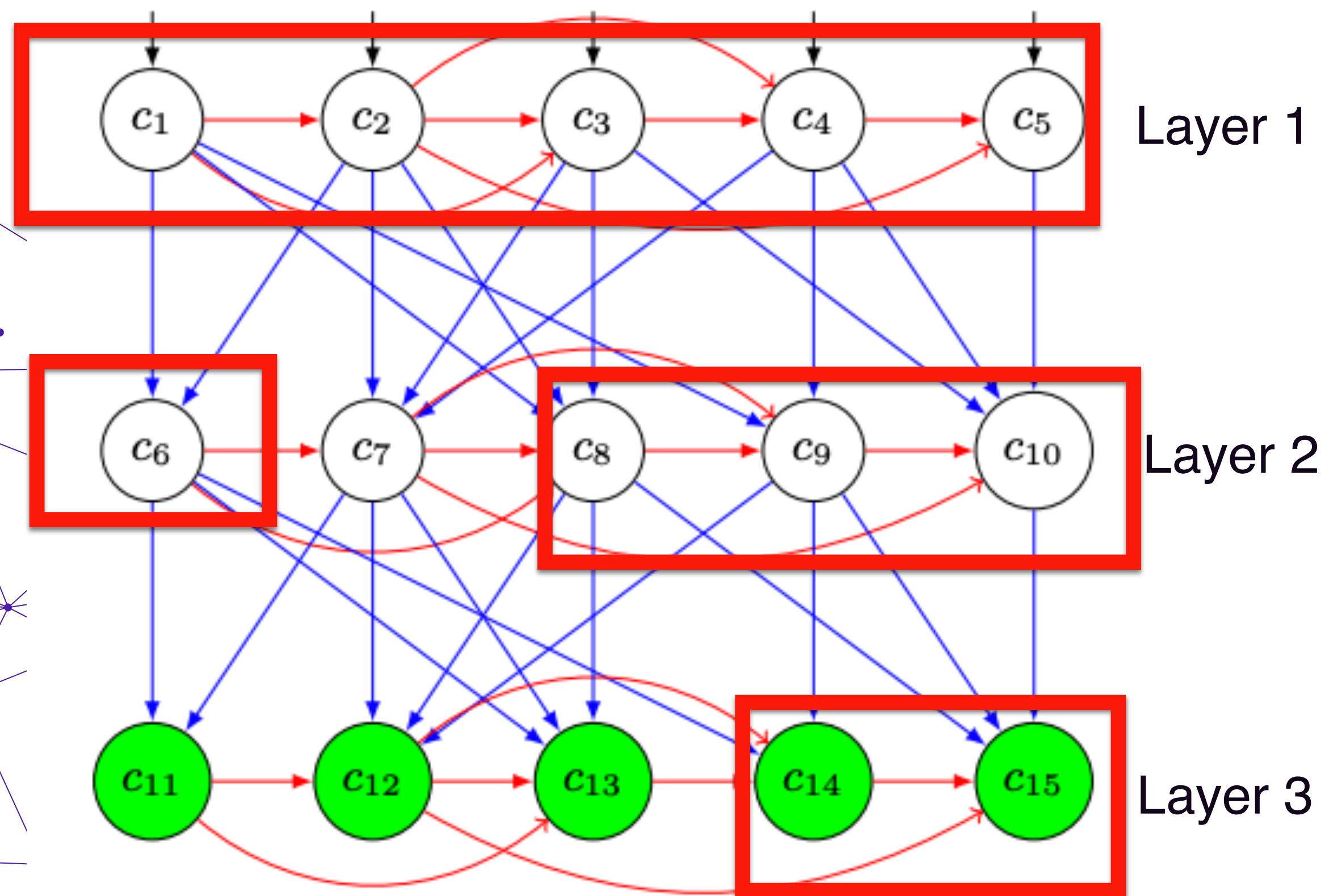


**Security argument:**

Recomputing c14, c15  
- 4 parents in layer 2 (because of the expander graph)

# Graph-labelling based PoS

## Stacked-DRGs graph (aka “SDR”)



### Security argument:

Recomputing  $c_{14}, c_{15}$

- 80% of the final layer => long path!!  
(because of the depth-robust graph)

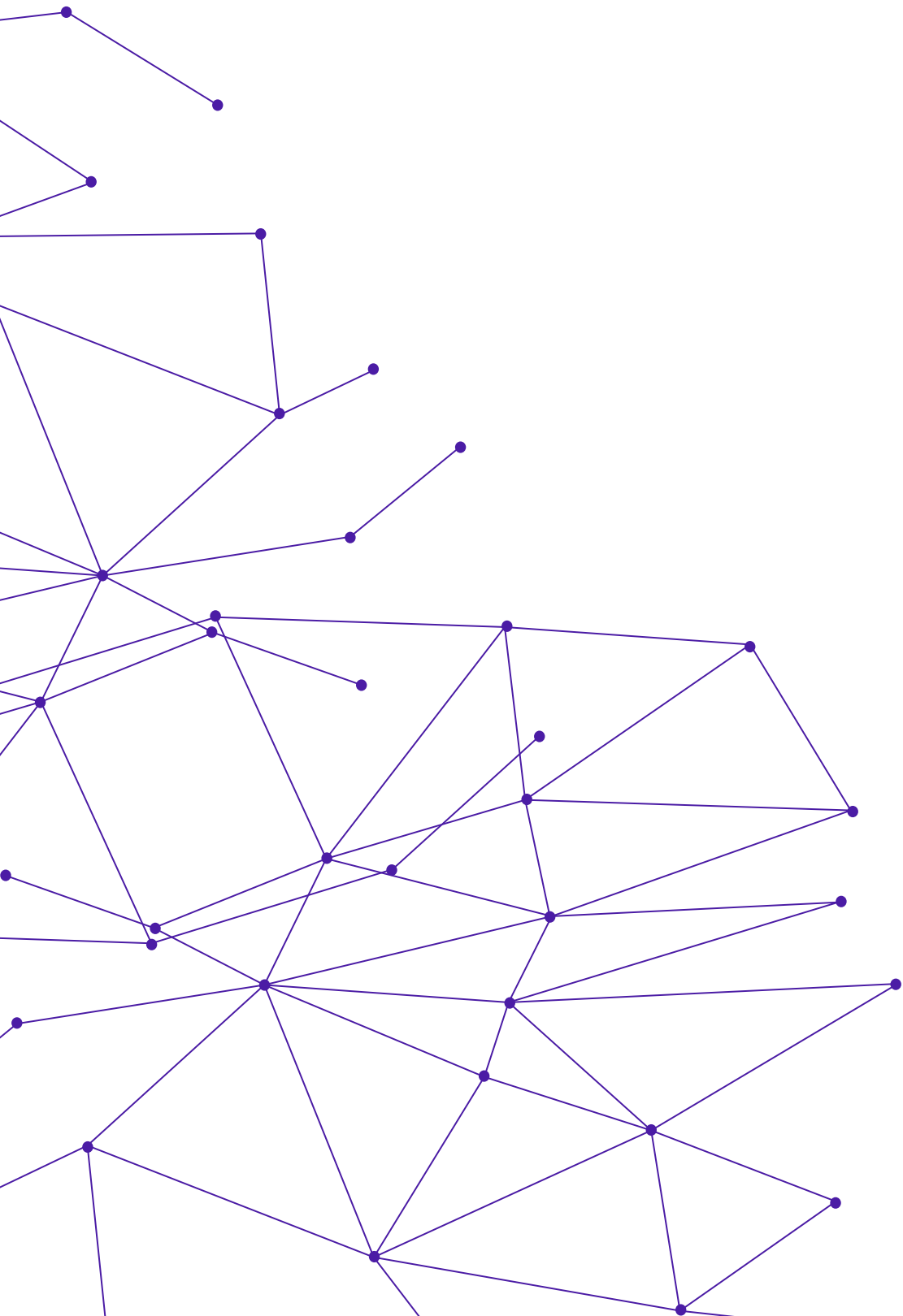
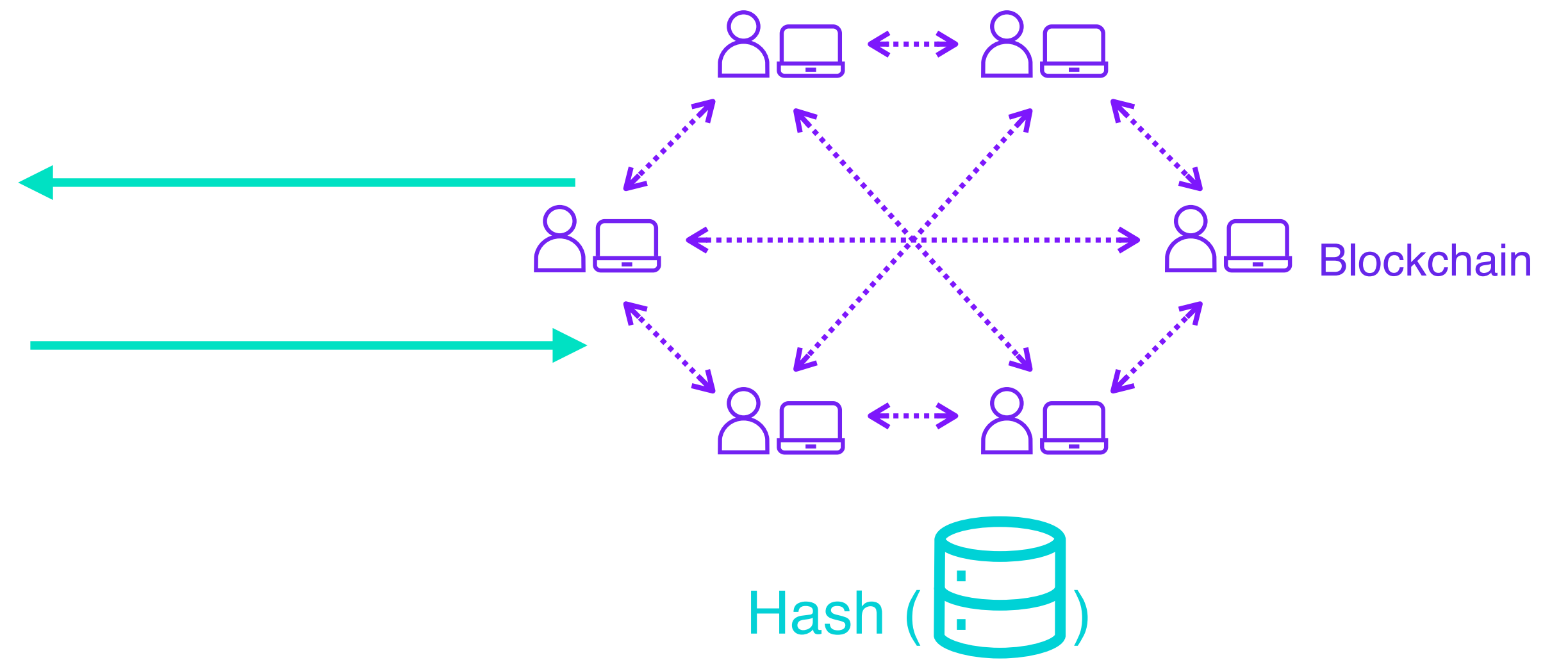
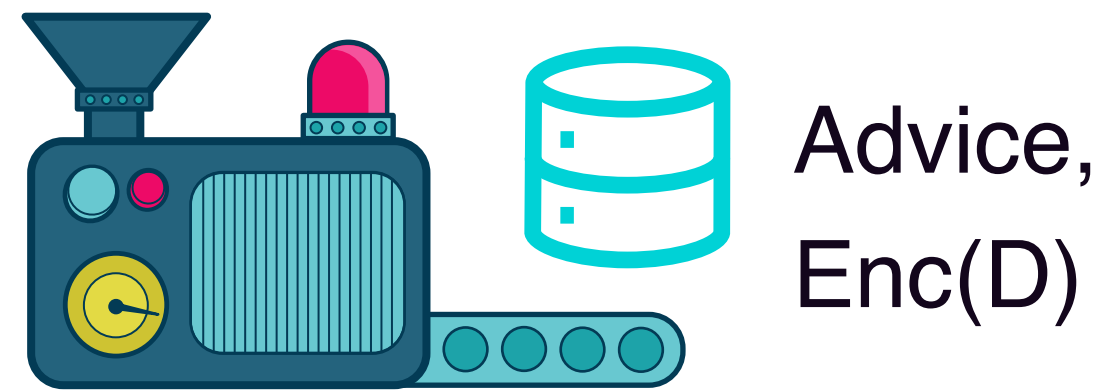
Long path => slow computation in the  
latency model

NOTE: We can also get a cost bound  
about the long path => cost model

# PoS in Filecoin



Storage Provider



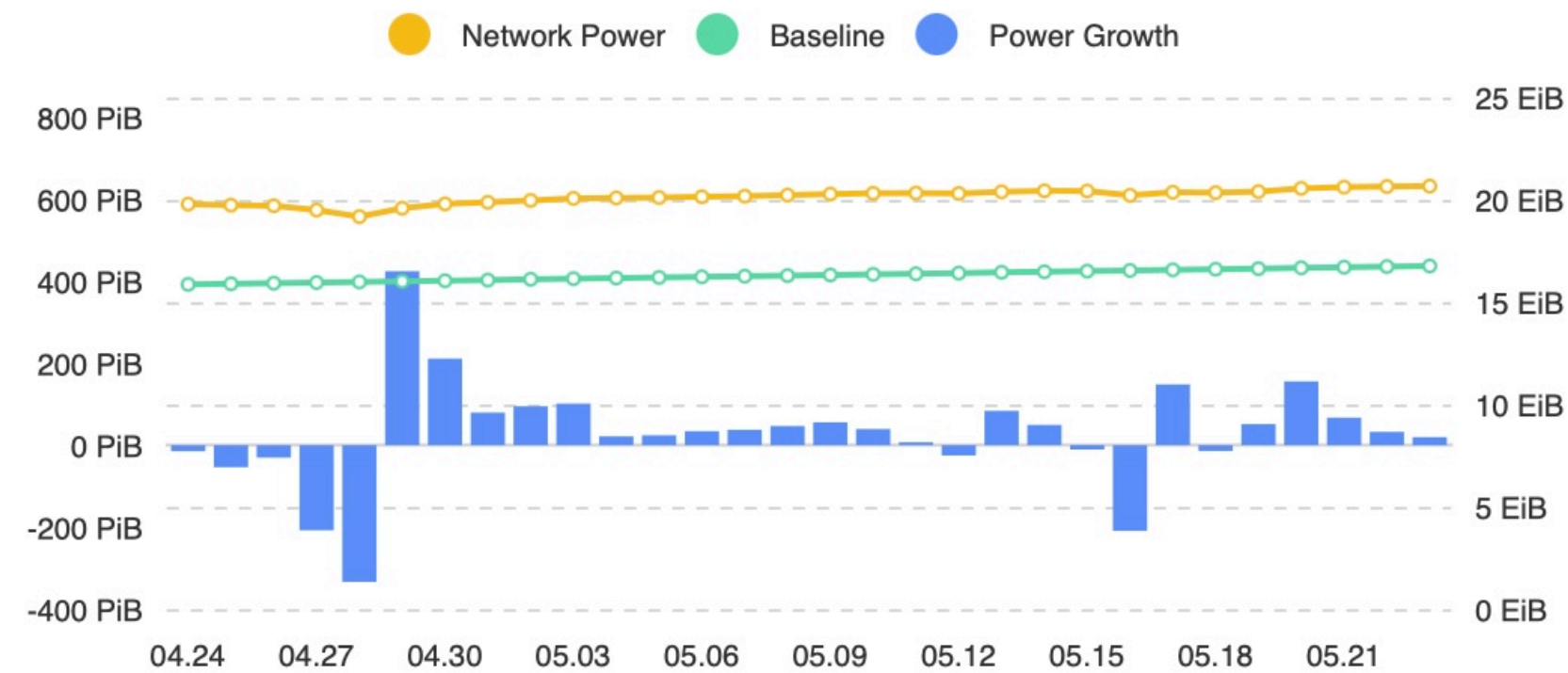
# PoS in Filecoin



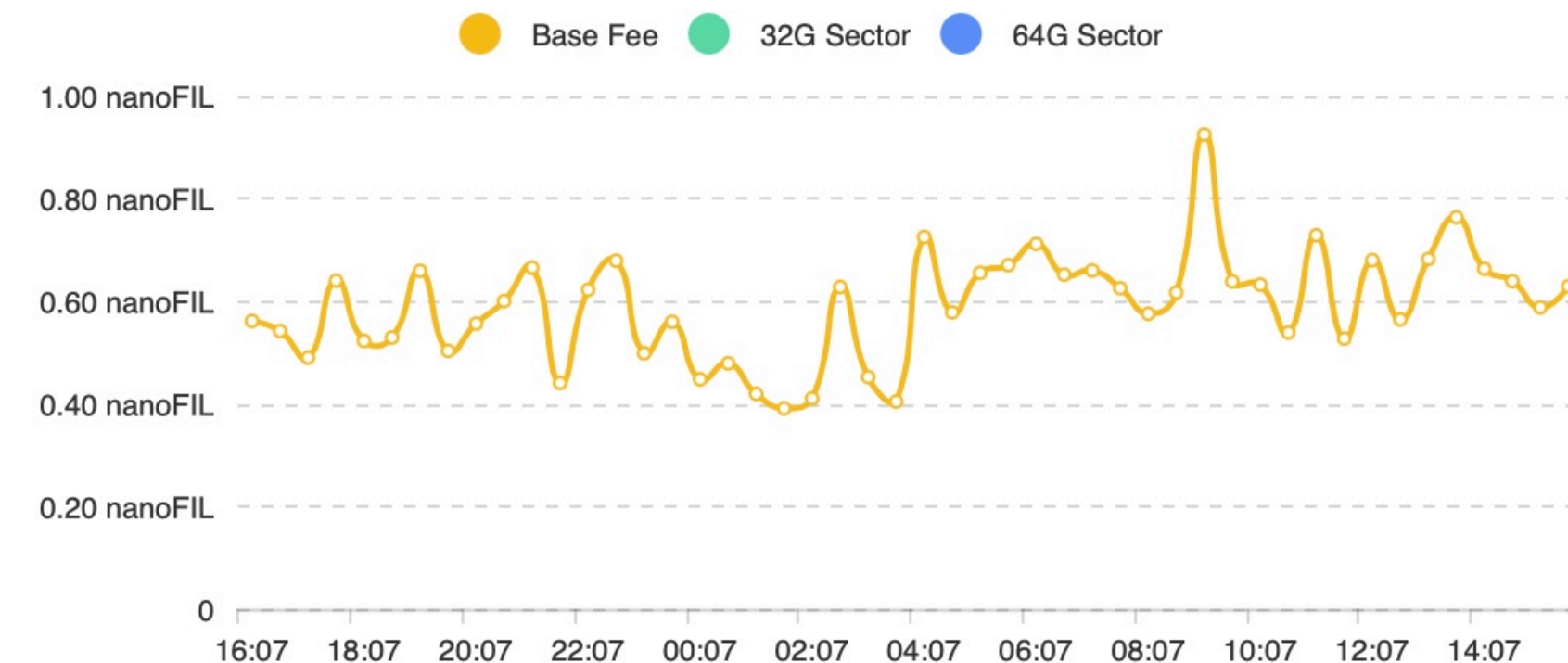
## Data index of the whole network open

Block Height <b>2,887,636</b>	Latest Block <b>2mins 14s</b>	Network Storage Power ⓘ <b>20.7046 EiB</b>	Latest 24h Power Growth <b>48.2998 PiB</b>	Latest 24h Output Efficiency ⓘ <b>0.0097 FIL/T</b>
Base Fee <b>0.6036 nanoFIL</b>	Current Sector Initial Pledge <b>7.13925 FIL/TiB</b>	Gas Used of 32G Sectors ⓘ <b>1.6215 FIL/TiB</b>	Cost of Sealing 32G Sectors ⓘ <b>8.76073 FIL/TiB</b>	Latest 24h Block Reward <b>218,237.36 FIL</b>

## baseline and storage power trend ⓘ More >



## 24h Base Fee Variations More >



# PoS in Filecoin



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# NFT STORAGE

Free Storage for NFTs

Free decentralized storage and bandwidth for NFTs on IPFS and Filecoin.

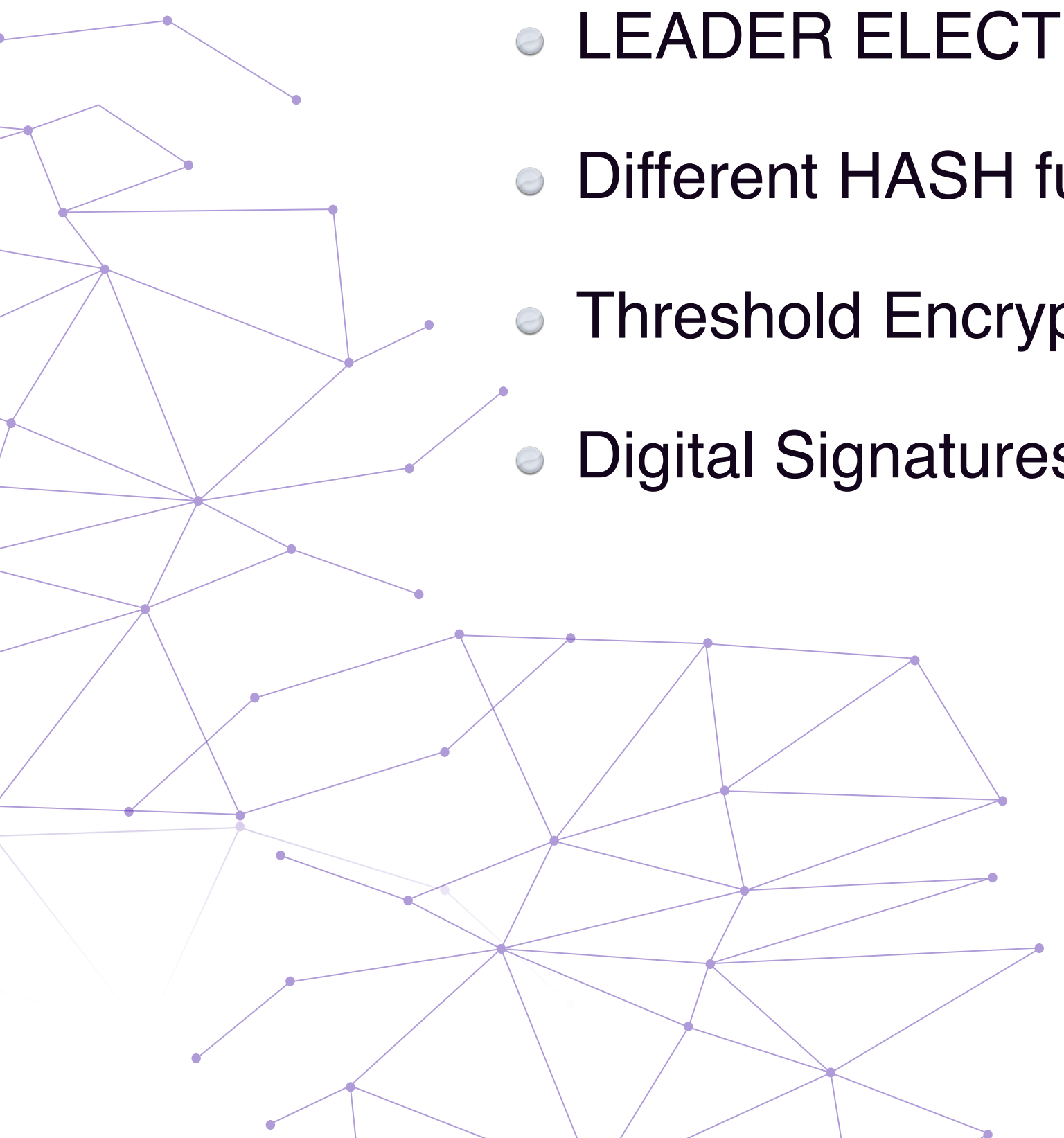
Get Started





## Is that all? Nope....

- SNARKs (“Succinct Non-interactive Argument of Knowledge”)
- VECTOR COMMITMENTS (ie, Merkle Trees)
- LEADER ELECTION via VRF (“verifiable random function”)
- Different HASH functions
- Threshold Encryption (for the random bacon, drand)
- Digital Signatures

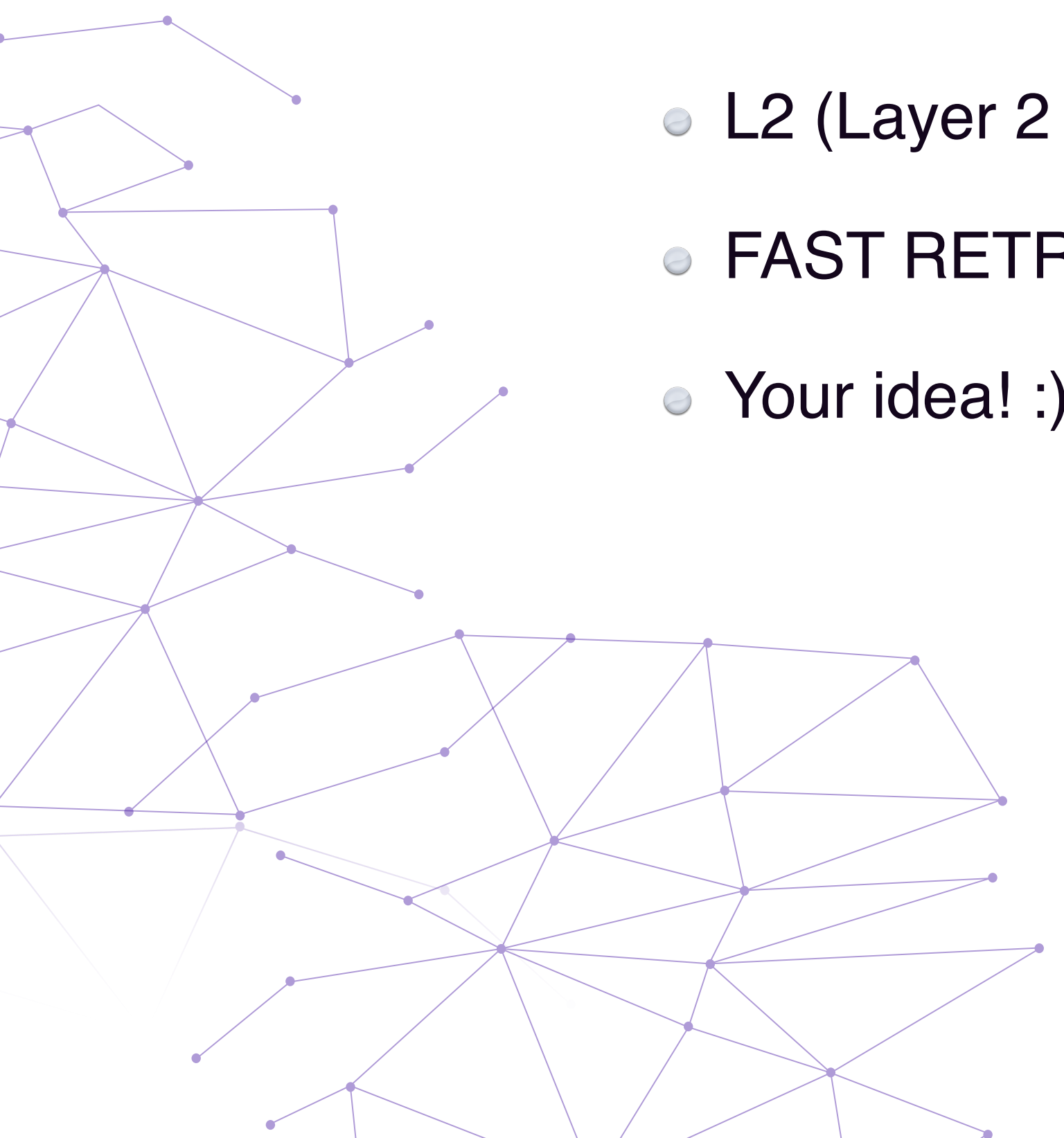


# PoS in Filecoin



## What's next? →

- FVM → Filecoin Virtual Machine and Smart Contract Capability
- L2 (Layer 2 Applications)
- FAST RETRIVABILITY
- Your idea! :)





# Thank you!

Questions?

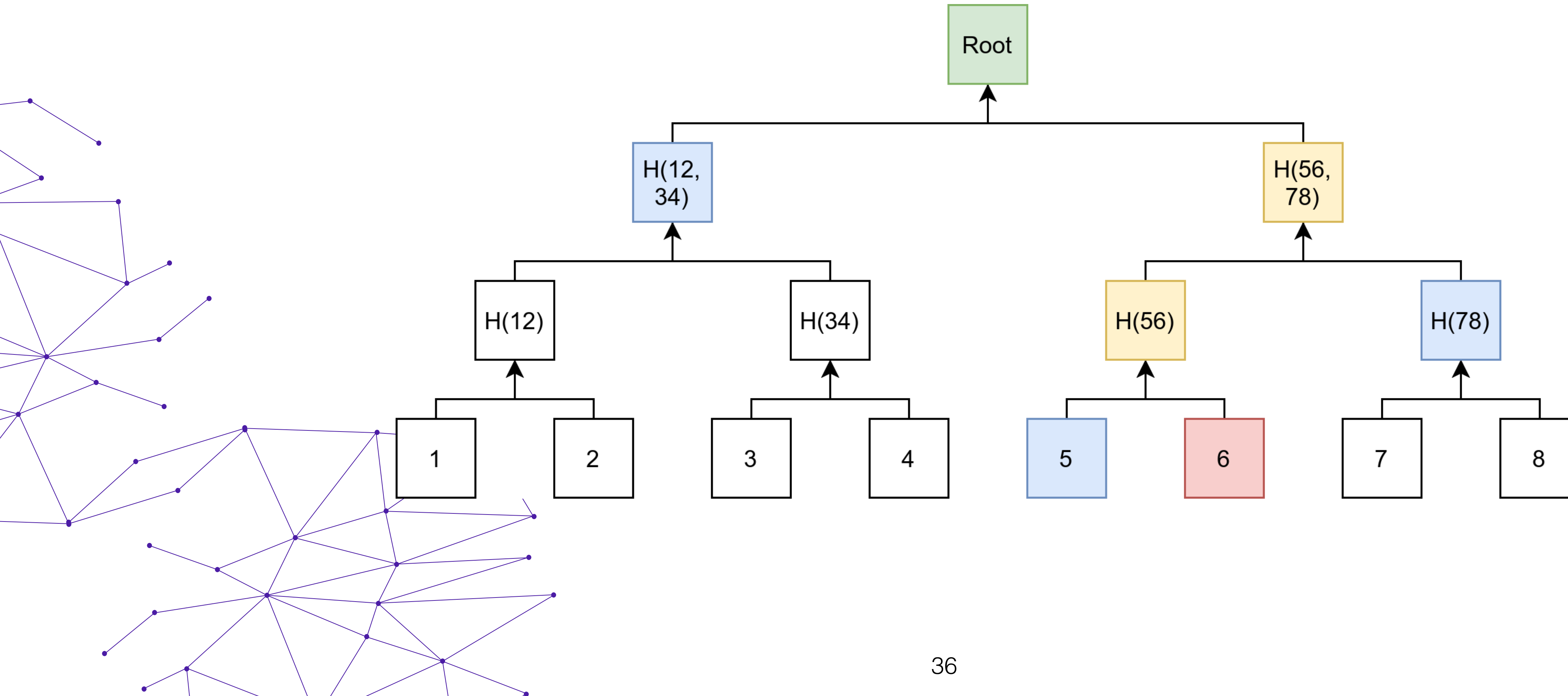
[irene.giacomelli@protocol.ai](mailto:irene.giacomelli@protocol.ai)  
[cryptonet.org](http://cryptonet.org)

# Recall: commitment scheme used in Filecoin



## How do we commit? Using Merkle-Trees (MT)

Root = Commitment (eg,  $\text{Com}(D)$  or  $\text{Com}(R)$ )

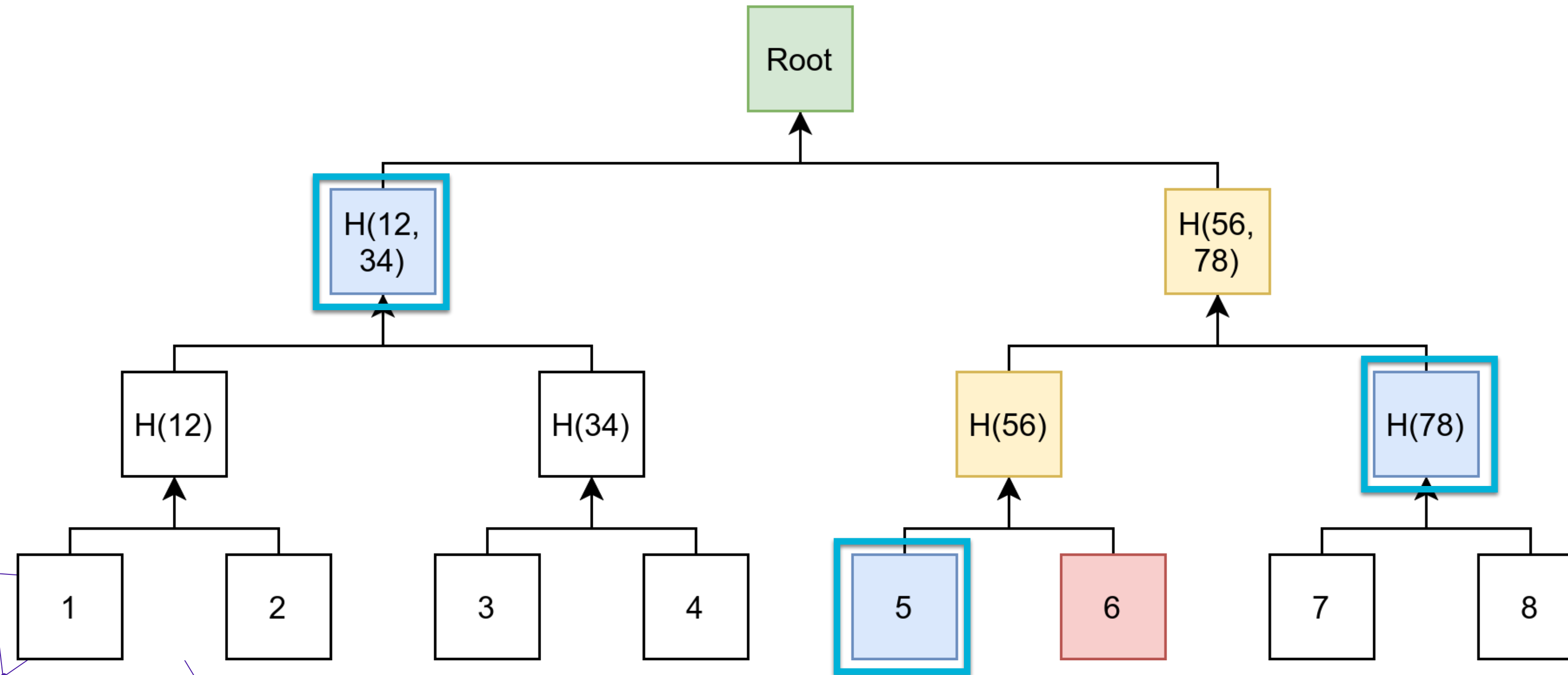


# Recall: commitment scheme used in Filecoin



## How do we commit? Using Merkle-Trees (MT)

Root = Commitment (eg, Com(D) or Com(R))



Blue values = proof that 6th value is correct respect to the commitment

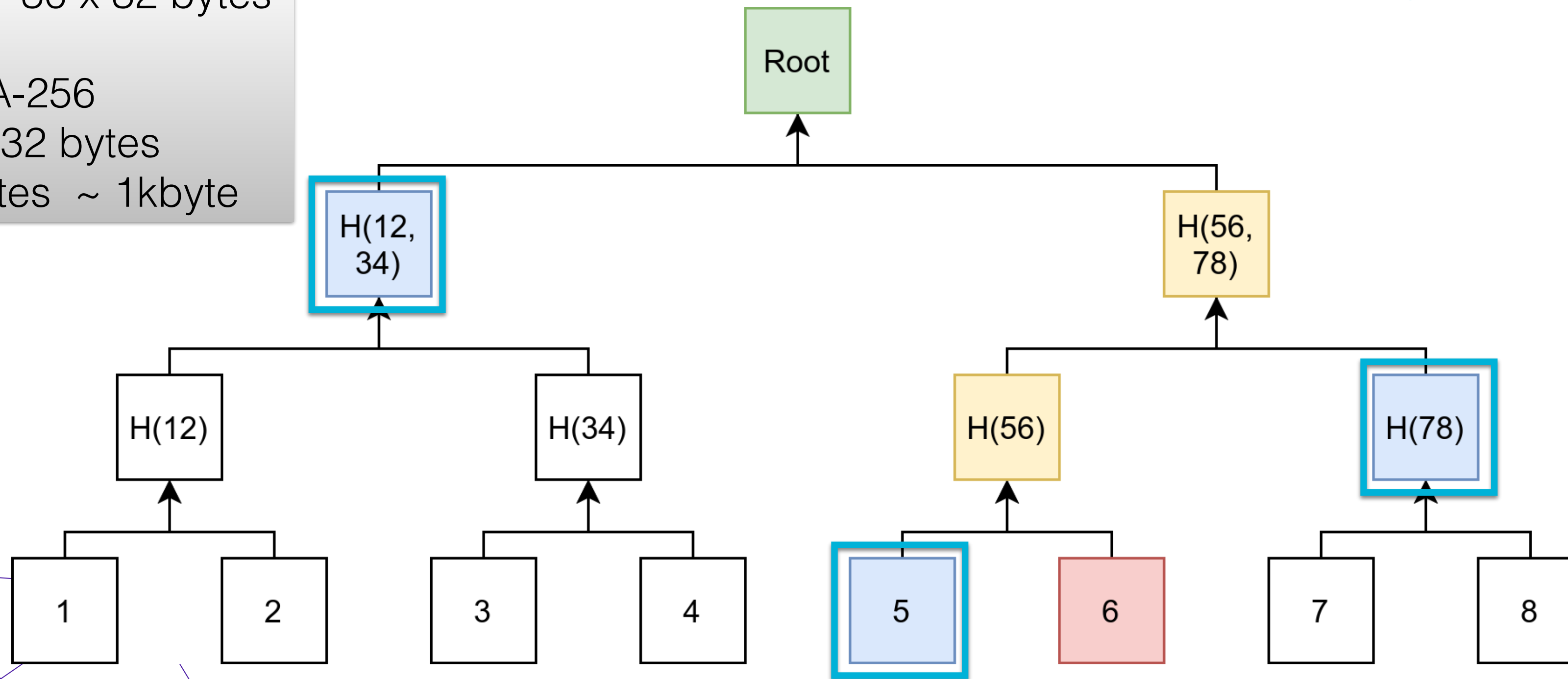
# Recall: commitment scheme used in Filecoin



## How do we commit? Using Merkle-Trees (MT)

Size of D = 32 GiB =  $2^{30} \times 32$  bytes  
Take H = SHA-256  
1 commitment = 32 bytes  
1 proof = 30 x 32 bytes ~ 1kbyte

Root = Commitment (eg, Com(D) or Com(R))



Blue values = proof that 6th value is correct respect to the commitment